

Data-Intensive Distributed Computing

CS 451/651 431/631 (Winter 2018)

Part 8: Analyzing Graphs, Redux (2/2)

March 22, 2018

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These slides are available at <http://lintool.github.io/bigdata-2018w/>



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Theme for Today:

How things work in the real world
(forget everything I told you...)

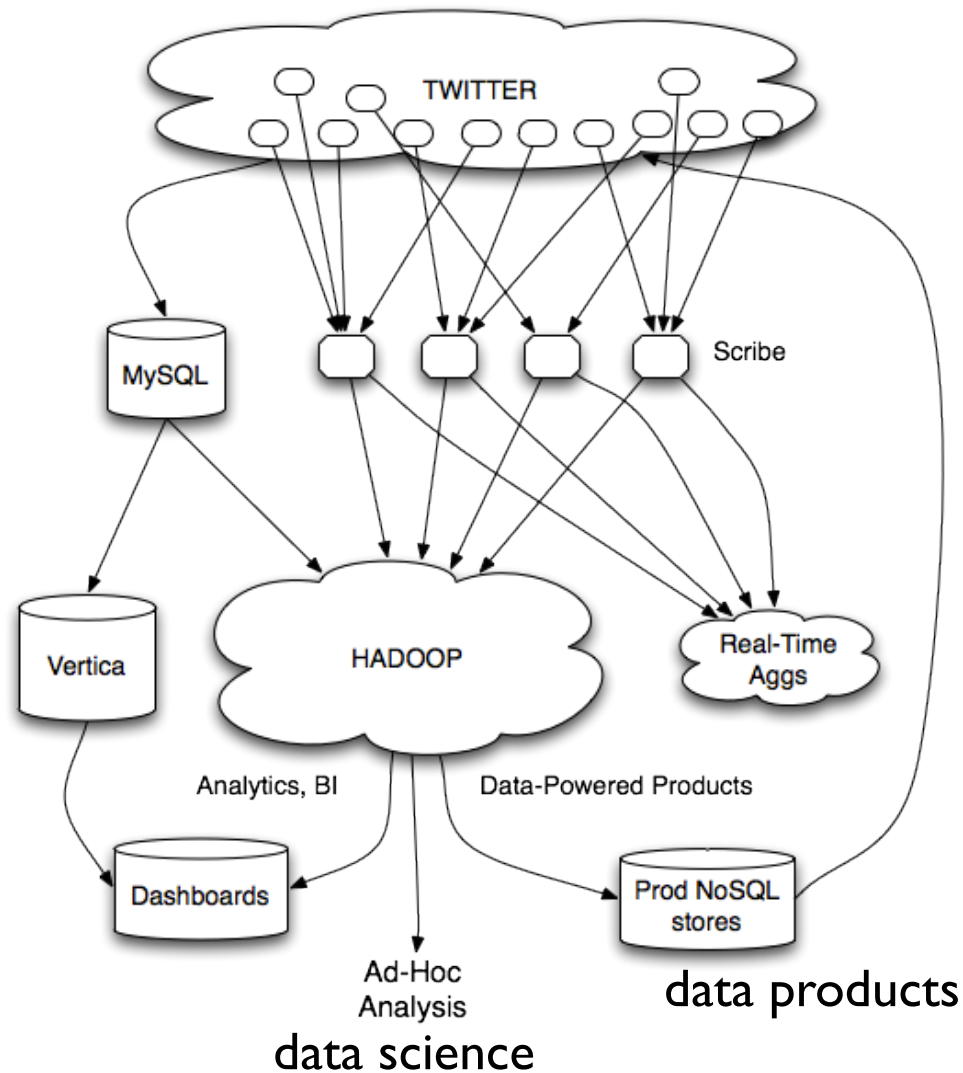


From the Ivory Tower...



... to building sh*t that works

What exactly did I do at Twitter?



I worked on...

- analytics infrastructure to support data science**
- data products to surface relevant content to users**

Mishne et al. Fast Data in the Era of Big Data: Twitter's Real-Time Related Query Suggestion Architecture. SIGMOD 2013.

@lintool

Tweets

TWEETS
1,641

Leibert et al. Automatic Management of Partitioned, Replicated Search Services. SoCC 2011

Compose new Tweet...

Who to follow · Refresh · View all



plotly @plotlygraphs
Follow Promoted



Brad Anderson @boorad
Followed by Florian Leibert ...
Follow



Sheila Morrissey @sheilaMorr
Follow

Popular accounts · Find friends

Trends

#Olymp
Ukraine
#Conf
Venny
#Premi

I worked on...

– analytics infrastructure to support data science

– data products to surface relevant content to users

cloudera
Struggling with complex data
of Data Science 2/20 to rethi
Promoted by Cloudera



Clinton Paquin @clintonpaquin
Simply stated, "The only pro
muscle memory" @TheChan
View conversation



The Hill @thehill · 1h
Republicans take debt ceiling
View summary



Retweeted by Alex Feinberg
Popehat @Popehat · 10h
In a world in which few thing
feed does.
Expand

sochi

#Sochi2014

#SochiProblems

Sochi

#SochiFail



Sochi 2014 @Sochi2014



Sochi Olympics 2014 @2014Sochi



Игры Сочи 2014 @sochi2014_ru



Sochi Problems @SochiProblem



NYT Olympics @SochiNYT



Sochi Problems @SochiProblems

Search all people for sochi

Gupta et al. WTF: The Who to Follow Service at Twitter. WWW 2013
Lin and Kolcz. Large-Scale Machine Learning at Twitter. SIGMOD 2012



circa ~2010

~150 people total

~60 Hadoop nodes

~6 people use analytics stack daily

circa ~2012

~1400 people total

10s of Ks of Hadoop nodes, multiple DCs

10s of PBs total Hadoop DW capacity

~100 TB ingest daily

dozens of teams use Hadoop daily

10s of Ks of Hadoop jobs daily




WTF

(who not to follow)

Who to follow · [refresh](#) · [view all](#)

-  **freshbooks** FreshBooks · [Follow](#)
 Promoted · Followed by @zappos and others. ×
-  **alanwarms** Alan Warms · [Follow](#)
Followed by @fredwilson and others. ×
-  **Mozzie21** Moises Henriques · [Follow](#)
can eat ×

Similar to @ryanhall3 · [view all](#)

-  **RunnerSpace_com** RunnerSpace.com · [Follow](#)
RunnerSpace.com has the latest in news and media...
-  **chrislieto** chris lieto · [Follow](#)
Chris Lieto is a top ranked World Class Triathlete, ...
-  **runningtimes** runningtimes · [Follow](#)

#numbers

(Second half of 2012)

~175 million active users

~20 billion edges

42% edges bidirectional

Avg shortest path length: 4.05

40% as many unfollows as follows daily

WTF responsible for ~1/8 of the edges

Graphs are core to Twitter



Graph-based recommendation systems
Why? Increase engagement!



The Journey

From the static follower graph for account recommendations...
... to the real-time interaction graph for content recommendations

In Four Acts...

In the beginning... the void

Act I
WTF and Cassovary

(circa 2010)

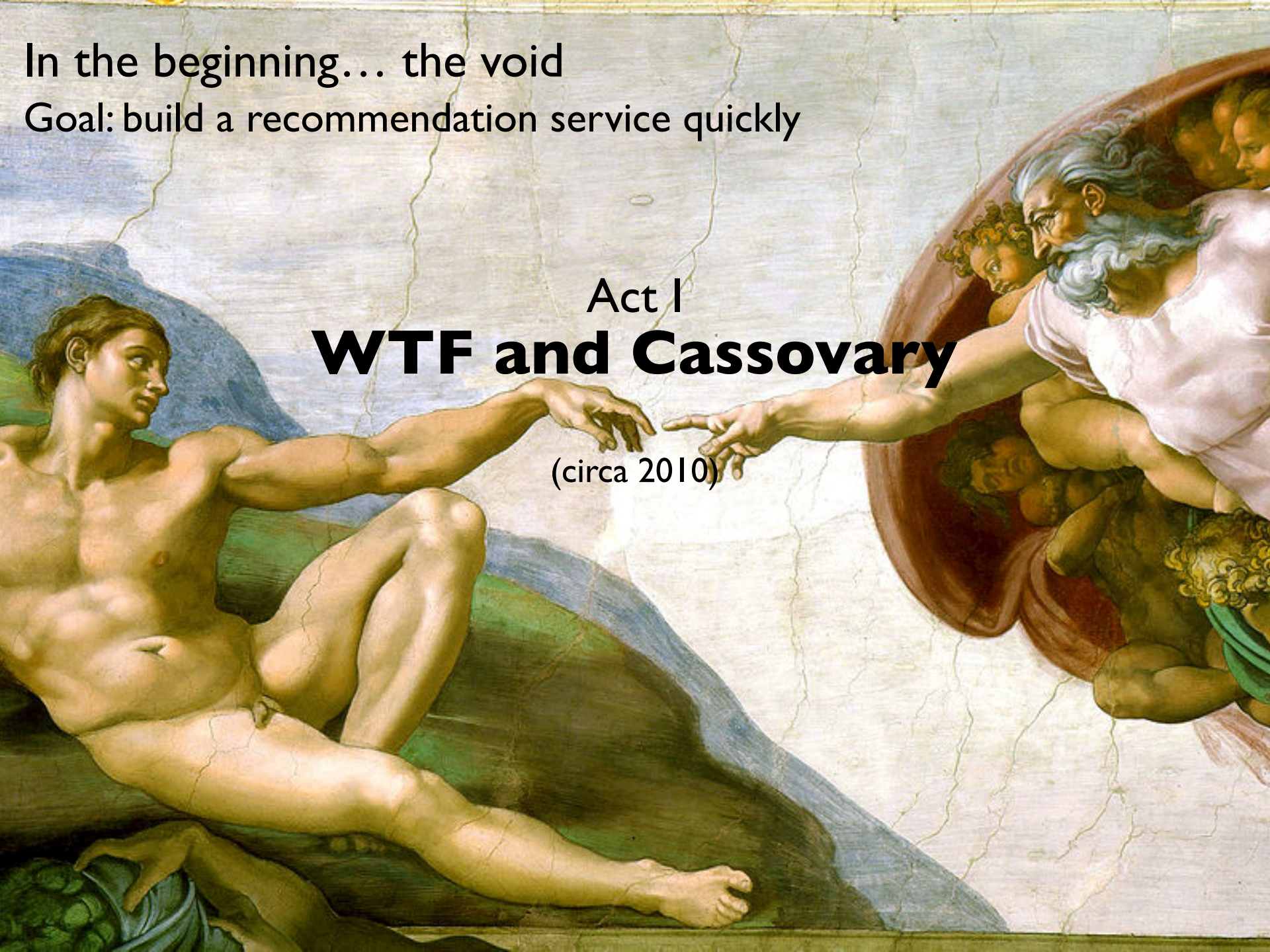
In the beginning... the void

Goal: build a recommendation service quickly

Act I

WTF and Cassovary

(circa 2010)



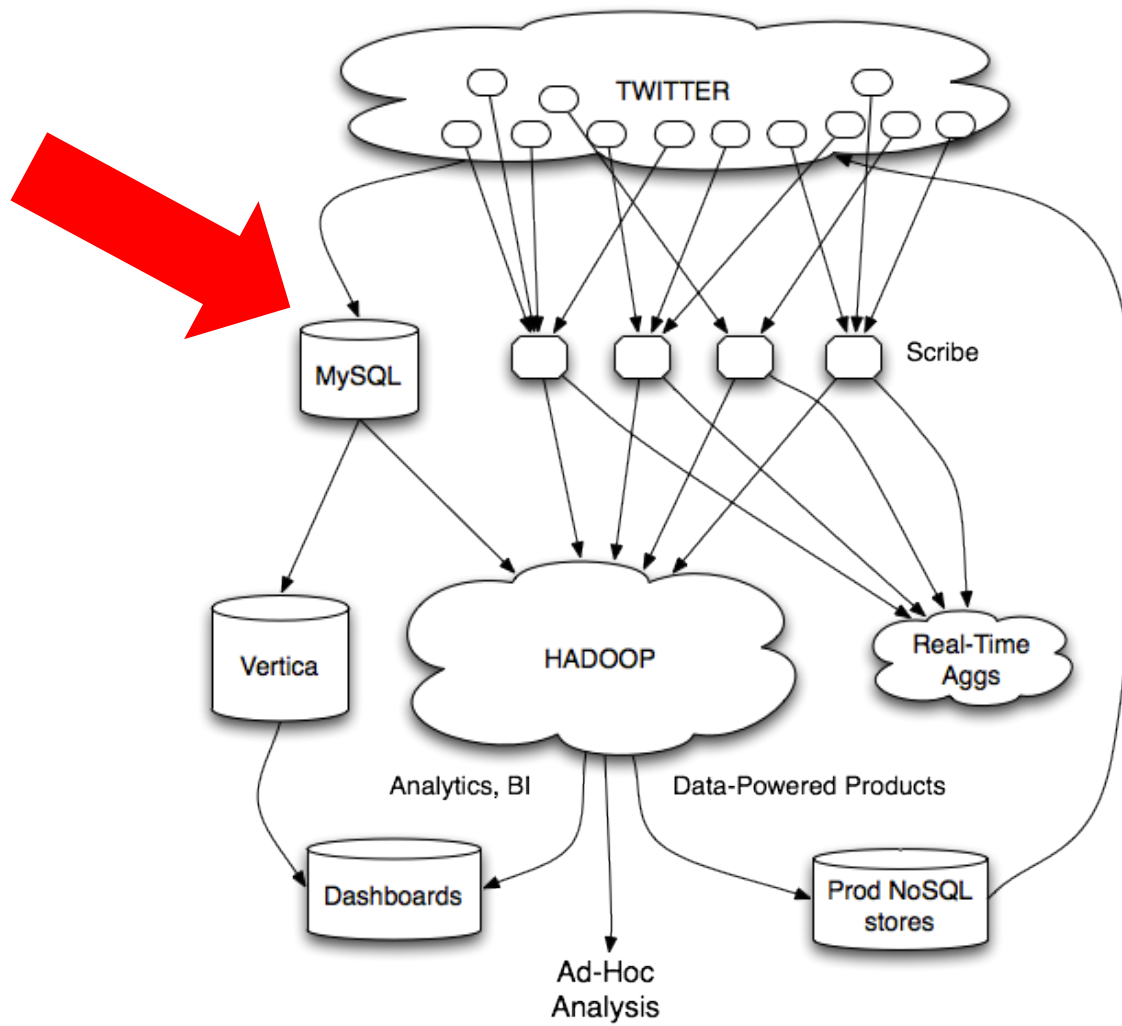


flockDB

(graph database)

Simple graph operations
Set intersection operations

Not appropriate for graph algorithms!



Okay, let's use MapReduce!
But MapReduce sucks for graphs!

What about...?

HaLoop (VLDB 2010)

Twister (MapReduce Workshop 2010)

Pregel/Giraph (SIGMOD 2010)

Graphlab (UAI 2010)

PrIter (SoCC 2011)

Datalog on Hyracks (Tech report, 2012)

Spark/GraphX (NSDI 2012, arXiv 2014)

PowerGraph (OSDI 2012)

GRACE (CIDR 2013)

Mizan (EuroSys 2013)

...

MapReduce sucks for graph algorithms...
Let's build our own system!

Key design decision:
Keep entire graph in memory... on a single machine!

Nuts!

Why?

Because we can!

Graph partitioning is hard... so don't do it

Simple architecture

Right choice at the time!

The runway argument



Suppose: 10×10^9 edges
(src, dest) pairs: ~ 80 GB

18×8 GB DIMMS = 144 GB

18×16 GB DIMMS = 288 GB

12×16 GB DIMMS = 192 GB

12×32 GB DIMMS = 384 GB

Cassovary

In-memory graph engine

- Implemented in Scala

- Compact in-memory representations

- But no compression

- Avoid JVM object overhead!

- Open-source



PageRank

“Semi-streaming” algorithm

Keep vertex state in memory, stream over edges

Each pass = one PageRank iteration

Bottlenecked by memory bandwidth

Convergence?

Don't run from scratch... use previous values

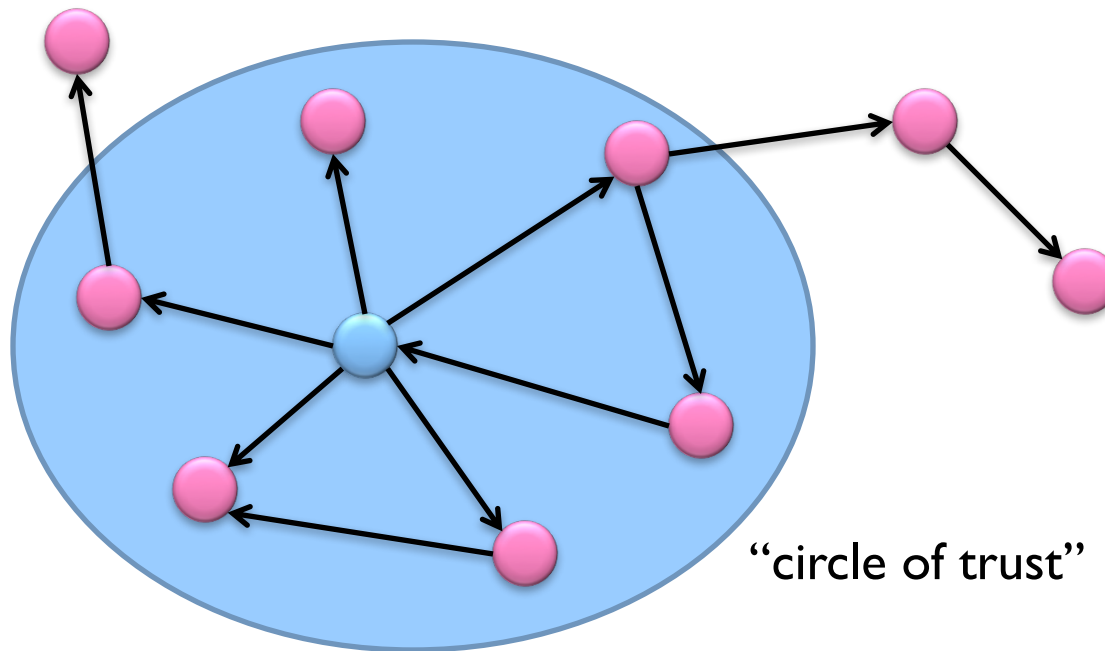
A few passes are sufficient

“Circle of Trust”

Ordered set of important neighbors for a user

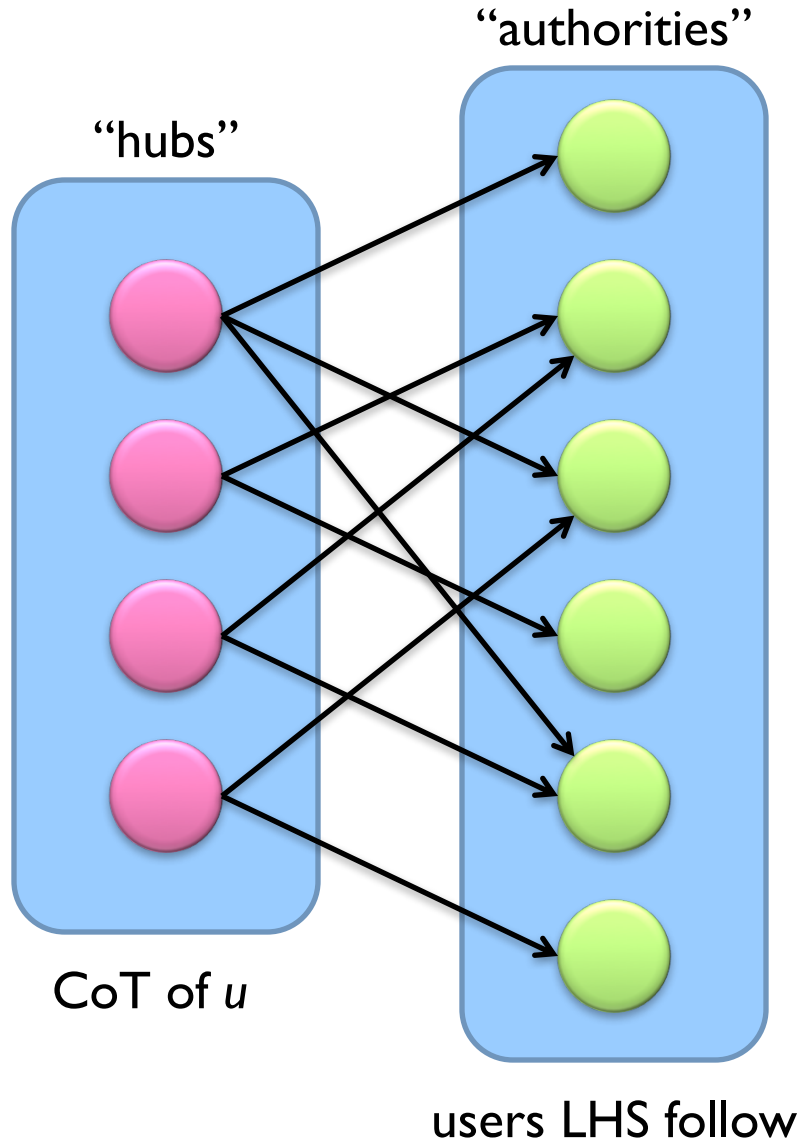
Result of egocentric random walk: Personalized PageRank!

Computed online based on various input parameters



One of the features used in search

SALSA for Recommendations

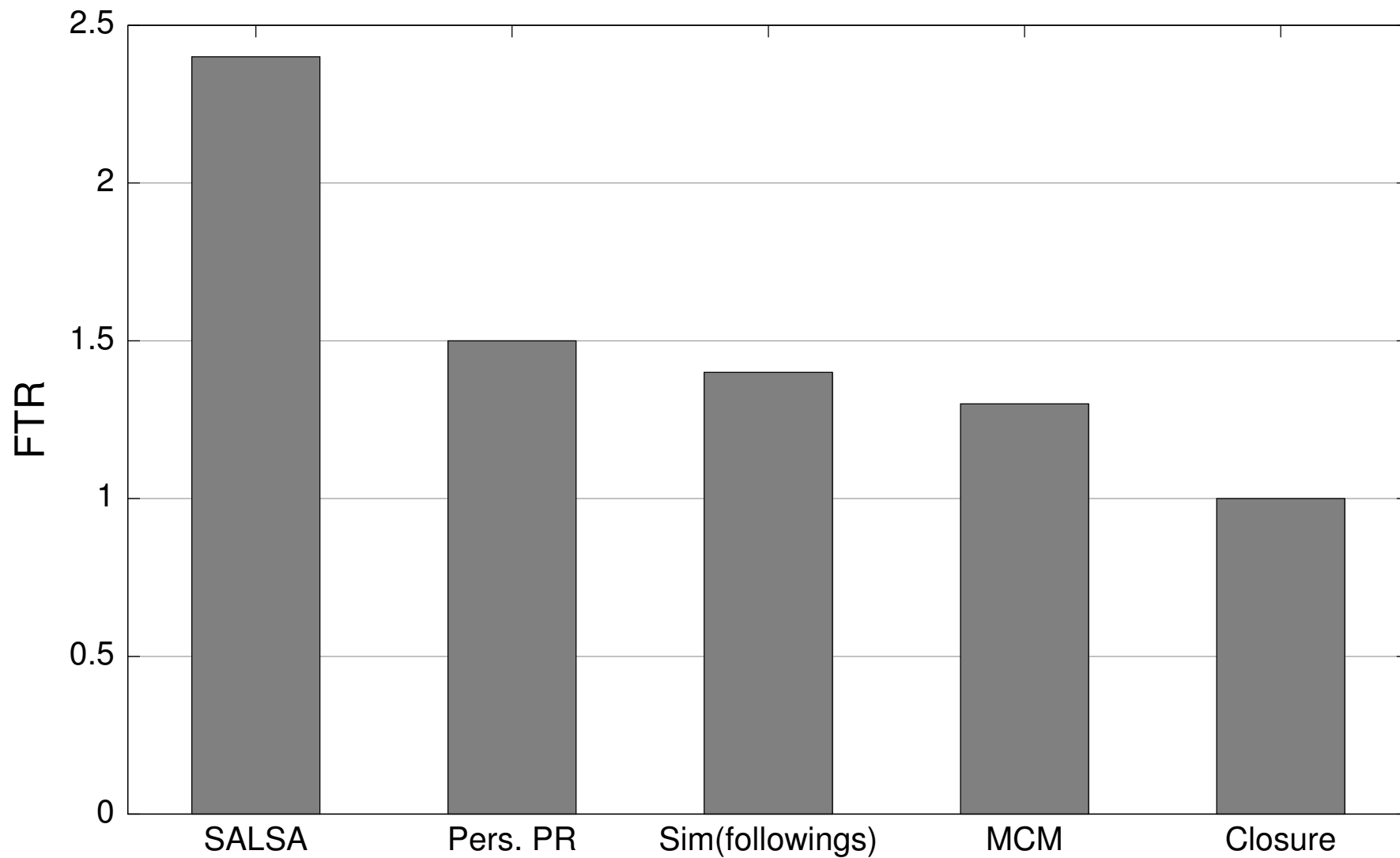


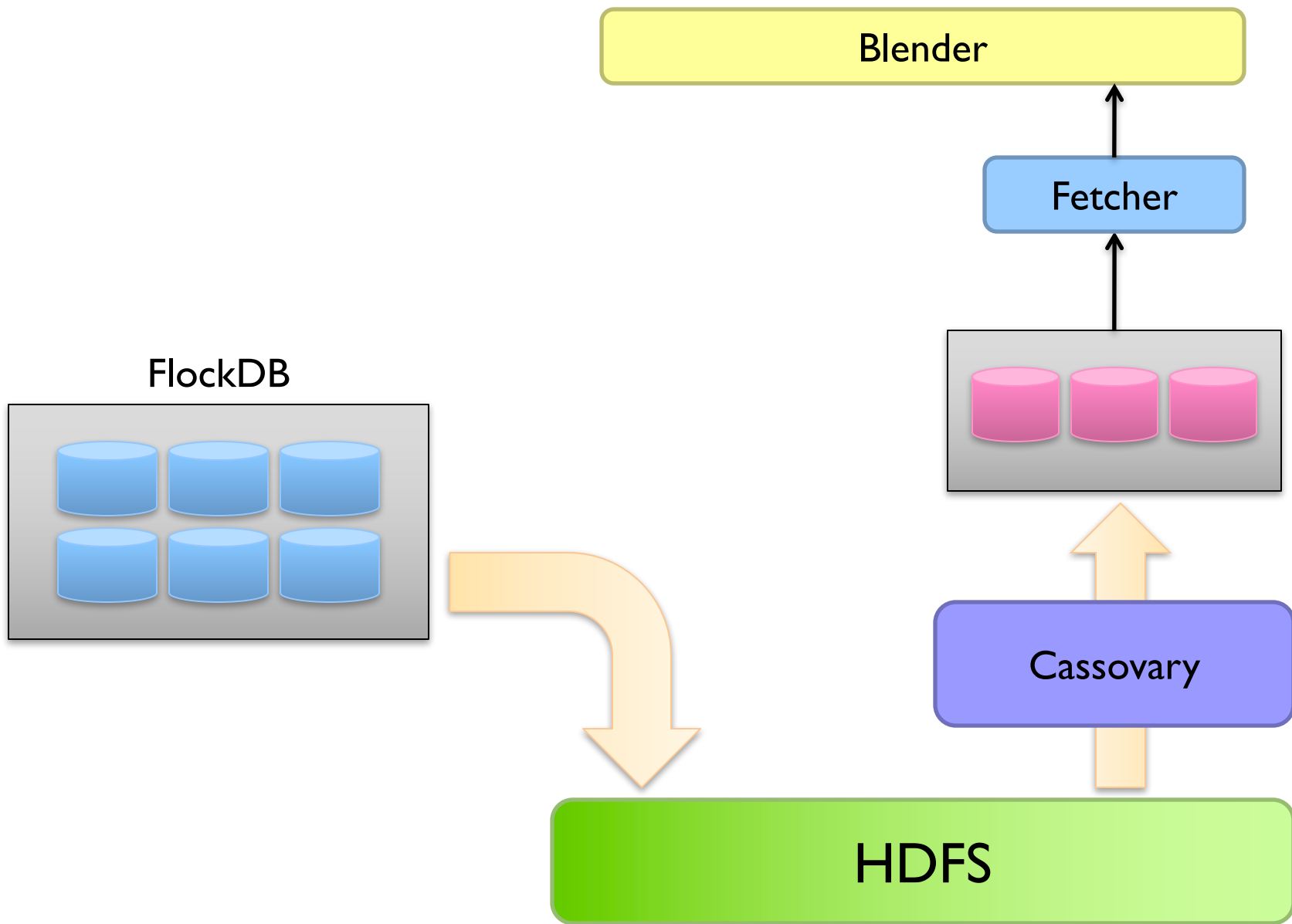
hubs scores:

similarity scores to u

authority scores:

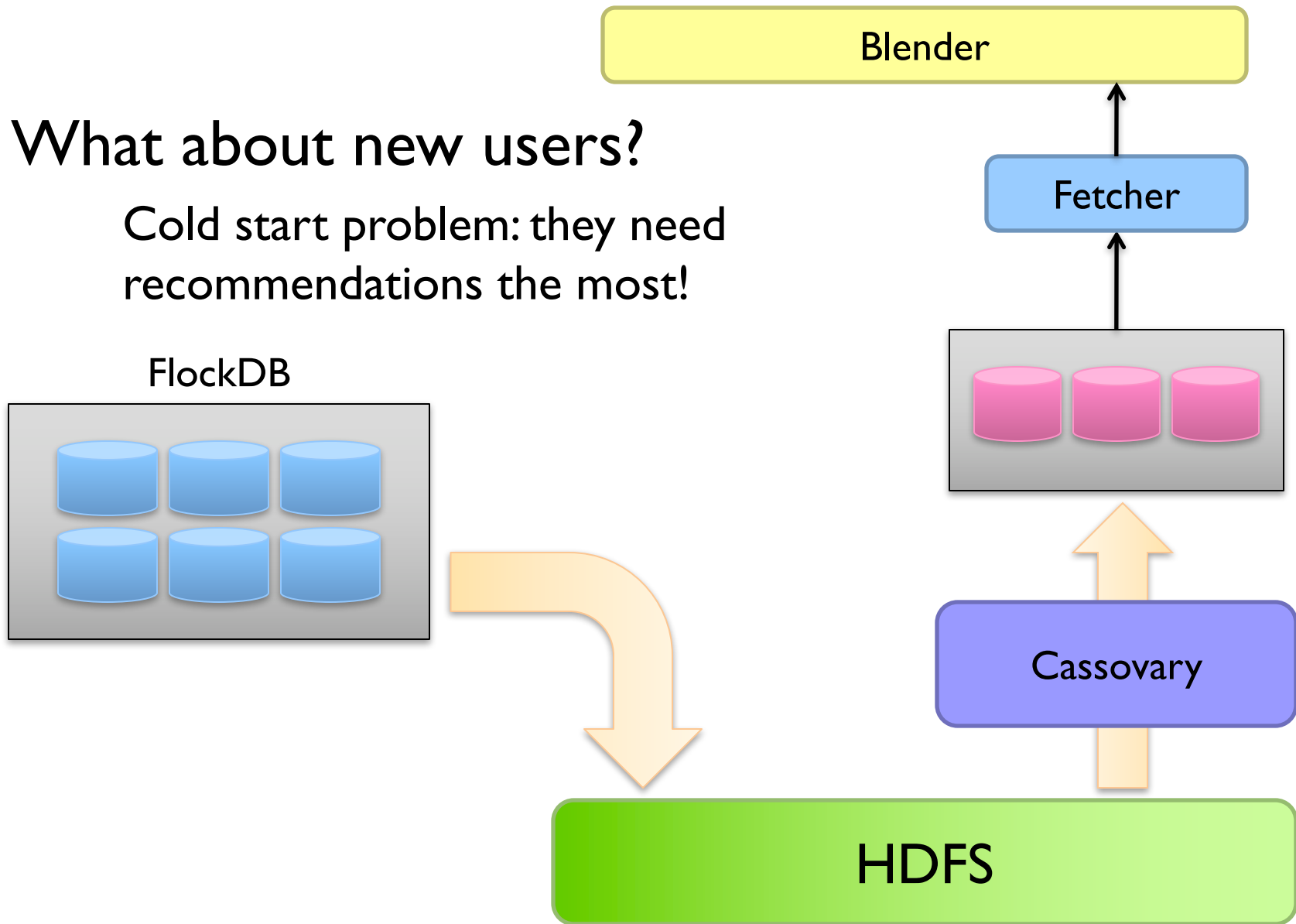
recommendation scores for u

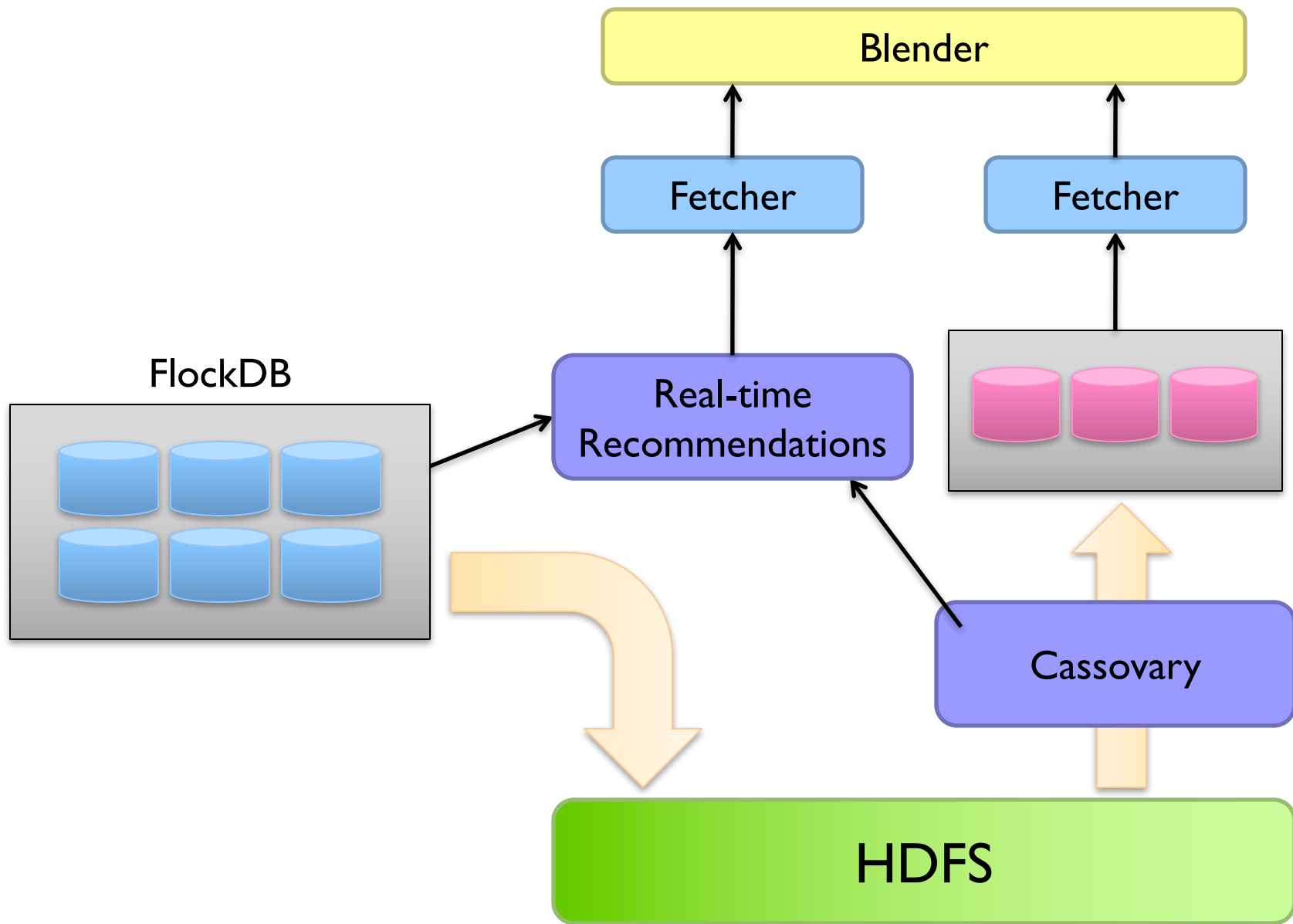




What about new users?

Cold start problem: they need recommendations the most!





Spring 2010: no WTF
seriously, WTF?

Summer 2010: WTF launched

A photograph of a server room with blue and red lighting. The room is filled with rows of server racks. The floor is dark, and the ceiling has exposed pipes and cables. The text "Act II RealGraph" is overlaid in the center, with "(circa 2012)" below it.

Act II **RealGraph**

(circa 2012)

Goel et al. Discovering Similar Users on Twitter. MLG 2013.



Another “interesting” design choice:
We migrated from Cassovary back to Hadoop!

Whaaaaaa?

Cassovary was a stopgap!

Hadoop provides:

Richer graph structure

Simplified production infrastructure

Scaling and fault-tolerance “for free”

Right choice at the time!

Wait, didn't you say MapReduce sucks?

What exactly is the issue?

Random walks on egocentric 2-hop neighborhood

Naïve approach: self-joins to materialize, then run algorithm

The shuffle is what kills you!

Graph algorithms in MapReduce

Tackle the shuffling problem!

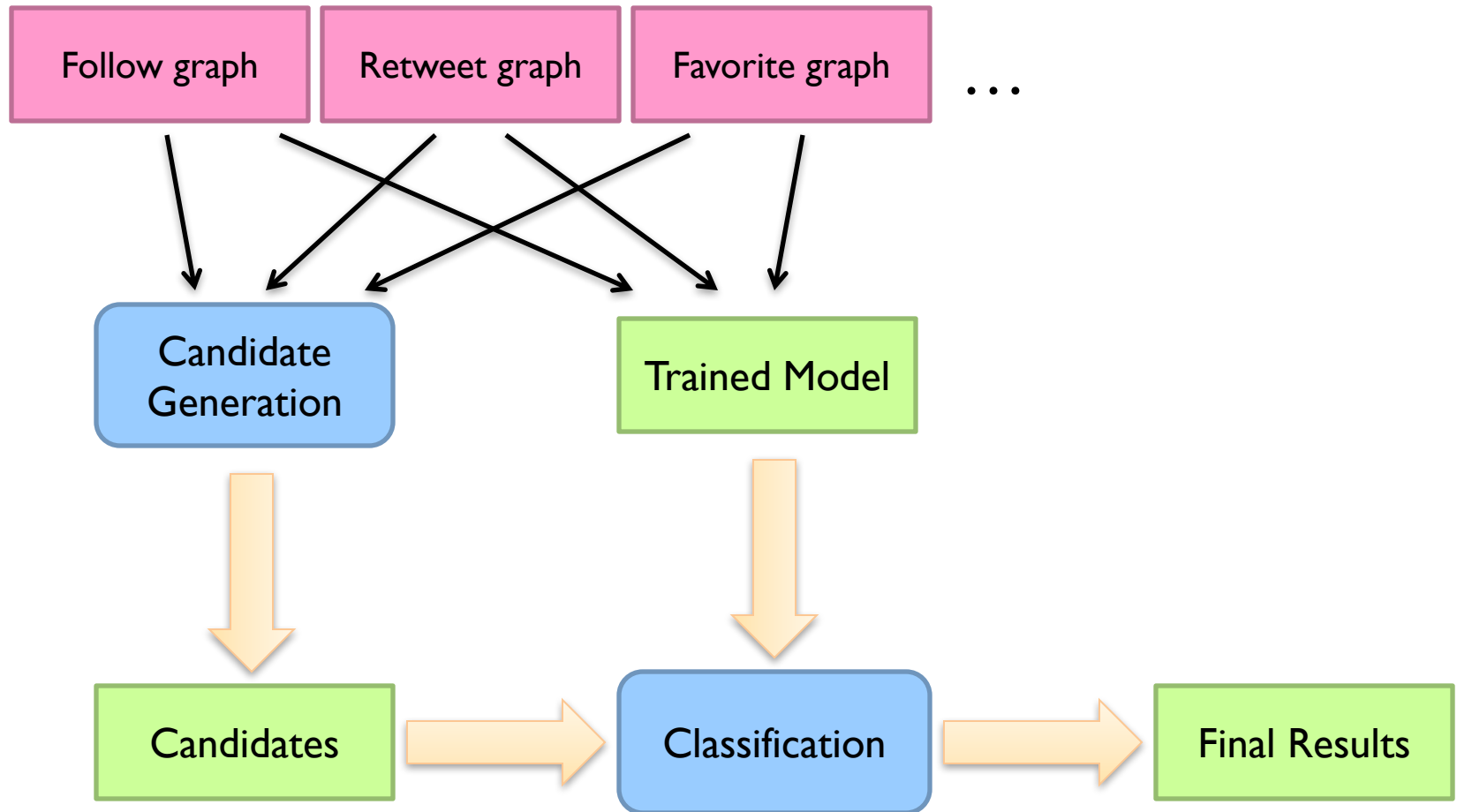
Key insights:

Batch and “stich together” partial random walks*

Clever sampling to avoid full materialization

* Sarma et al. Estimating PageRank on Graph Streams. PODS 2008
Bahmani et al. Fast Personalized PageRank on MapReduce. SIGMOD 2011.

Throw in ML while we're at it...





Act III
MagicRecs

(circa 2013)



@dickc
dick costolo

Our mission: Instantly connect people everywhere to what's most meaningful to them. #mwc11

14 Feb via web ☆ Favorite ↻ Retweet ↩ Reply

Isn't the point of Twitter real-time?

So why is WTF still dominated by batch processing?

Observation: fresh recommendations get better engagement

Logical conclusion: generate recommendations in real time!

From batch to real-time recommendations:

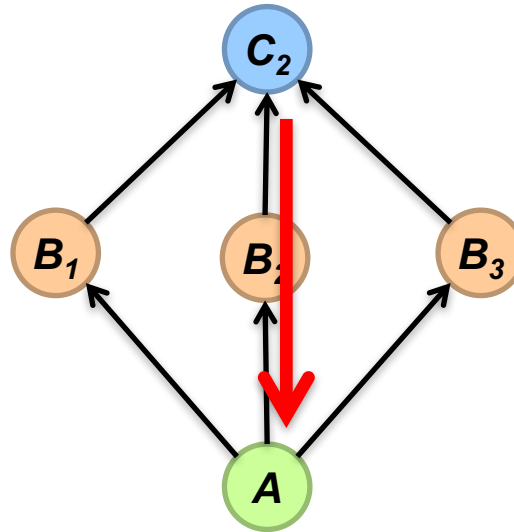
Recommendations based on recent activity

“Trending in your network”

Inverts the WTF problem:

For this user, what recommendations to generate?

Given this new edge, which user to make recommendations to?



Why does this work?

A follows B's because they're interesting
B's following C's because “something's happening”
(generalizes to any activity)

Scale of the Problem

$O(10^8)$ vertices, $O(10^{10})$ edges

Designed for $O(10^4)$ events per second

Naïve solutions:

Poll each vertex periodically

Materialize everyone's two-hop neighborhood, intersect

Production solution:

Idea #1: Convert problem into adjacency list intersection

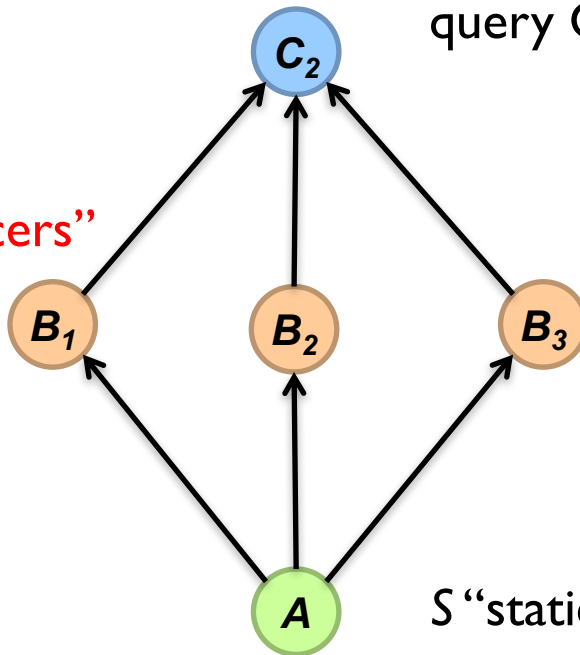
Idea #2: Partition graph to eliminate non-local intersections

Single Node Solution

D “dynamic” structure:
stores inverted adjacency lists
query *C*, return all *B*’s that link to it

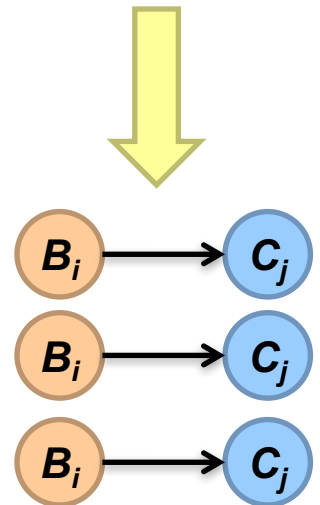
Who we’re recommending

“influencers”



Who we’re making the
recommendations to

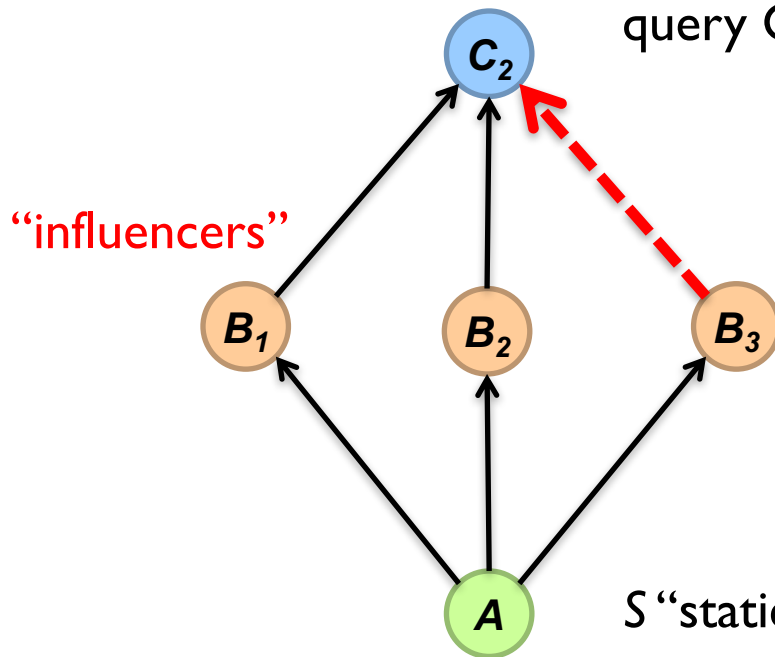
S “static” structure:
stores inverted adjacency lists
query *B*, return all *A*’s that link to it



Algorithm

D “dynamic” structure:
stores inverted adjacency lists
query C, return all B’s that link to it

Who we’re recommending

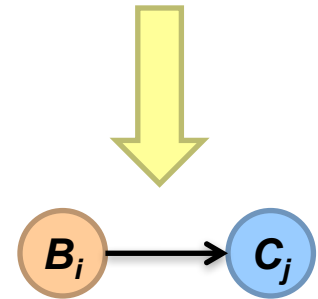


1. Receive B_3 to C_2
2. Query D for C_2 , get B_1, B_2, B_3
3. For each B_1, B_2, B_3 , query S
4. Intersect lists to compute A 's

S “static” structure:
stores inverted adjacency lists
query B, return all A’s that link to it

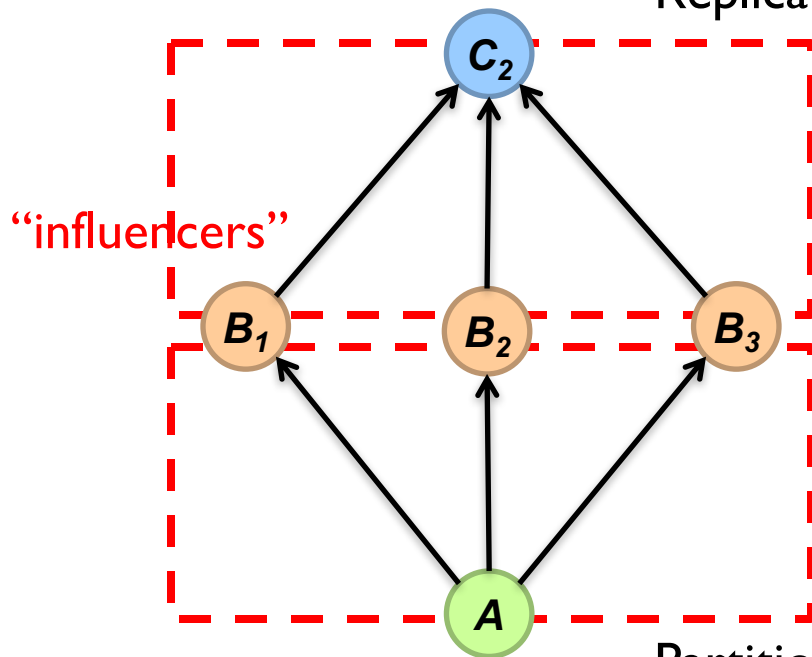
Idea #1: Convert problem into adjacency list intersection

Distributed Solution

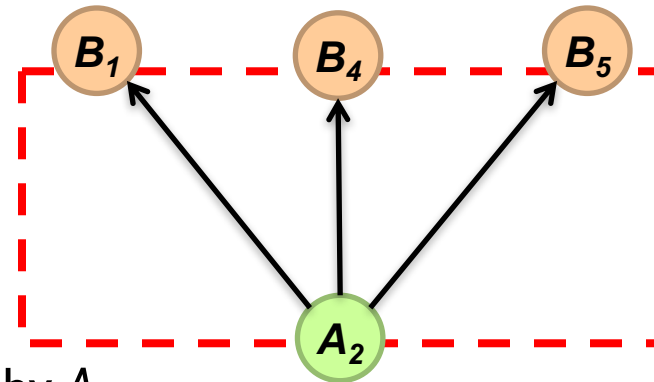


Who we're recommending

Replicate on every node



1. Fan out new edge to every node
2. Run algorithm on each partition
3. Gather results from each partition



Partition by A

Who we're making the recommendations to

Idea #2: Partition graph to eliminate non-local intersections

Production Status

Launched September 2013

Usage Statistics (Circa 2014)

Push recommendations to Twitter mobile users

Billions of raw candidates, millions of push notifications daily

Performance

End-to-end latency (from edge creation to delivery):

median 7s, p99 15s



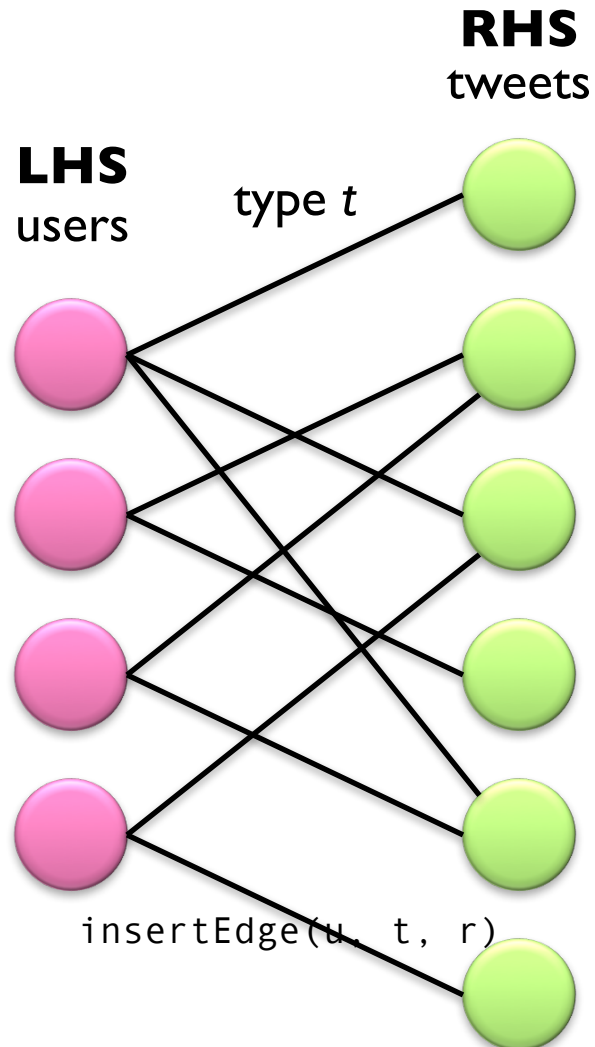
Act IV
GraphJet

(circa 2014)

Fully bought into the potential of real-time...
but needed something more general

Focused specifically on the interaction graph

Data Model



`getLeftVertexEdges(u)`
`getLeftVertexRandomEdges(u, k)`

`getRightVertexEdges(t)`
`getRightVertexRandomEdges(t, k)`

Noteworthy design decisions

Make it simple, make it fast!

No partitioning

Focus on recent data, fits on a single machine

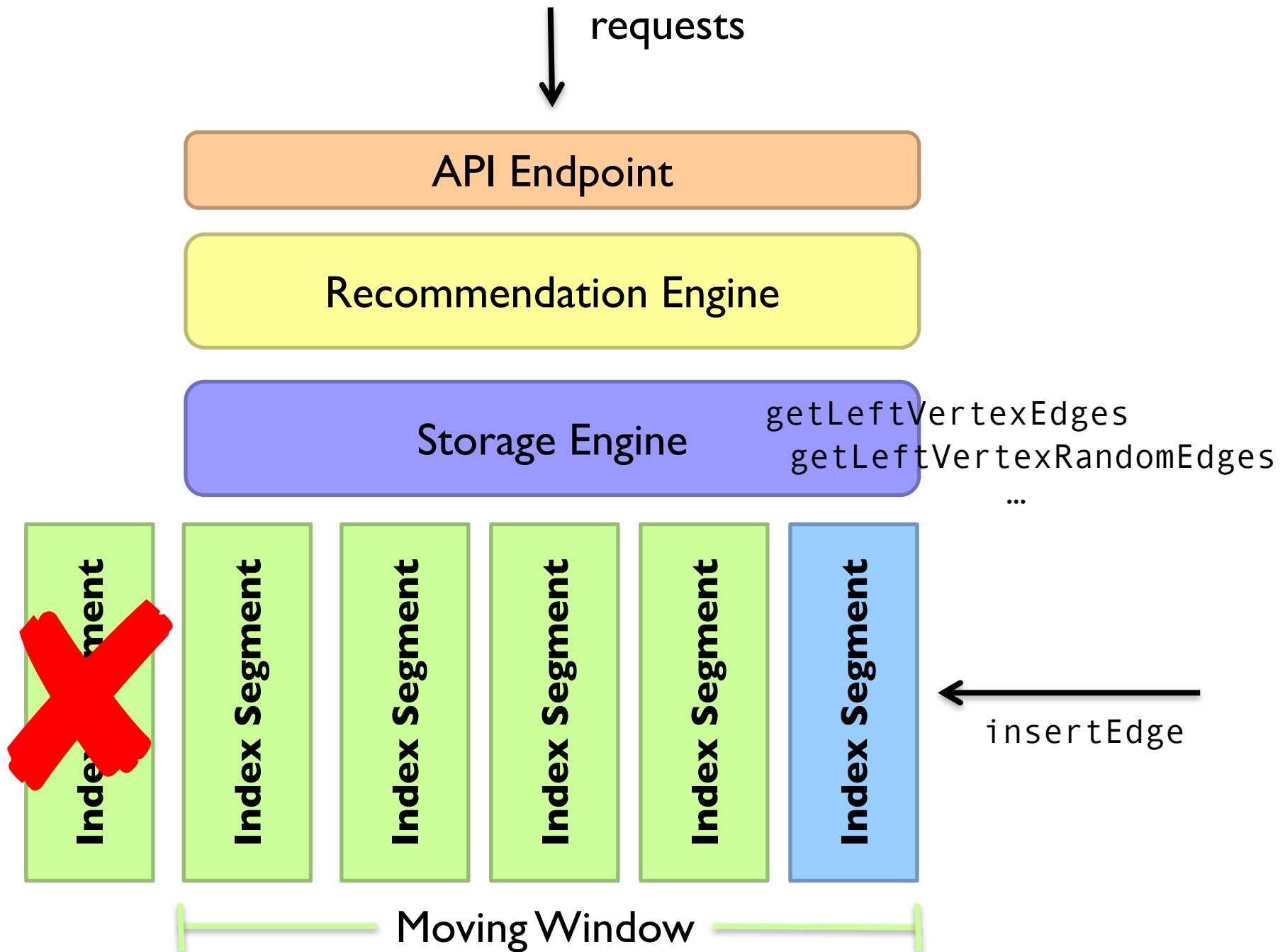
No deletes

Not meaningful w/ interaction data

No arbitrary edge metadata

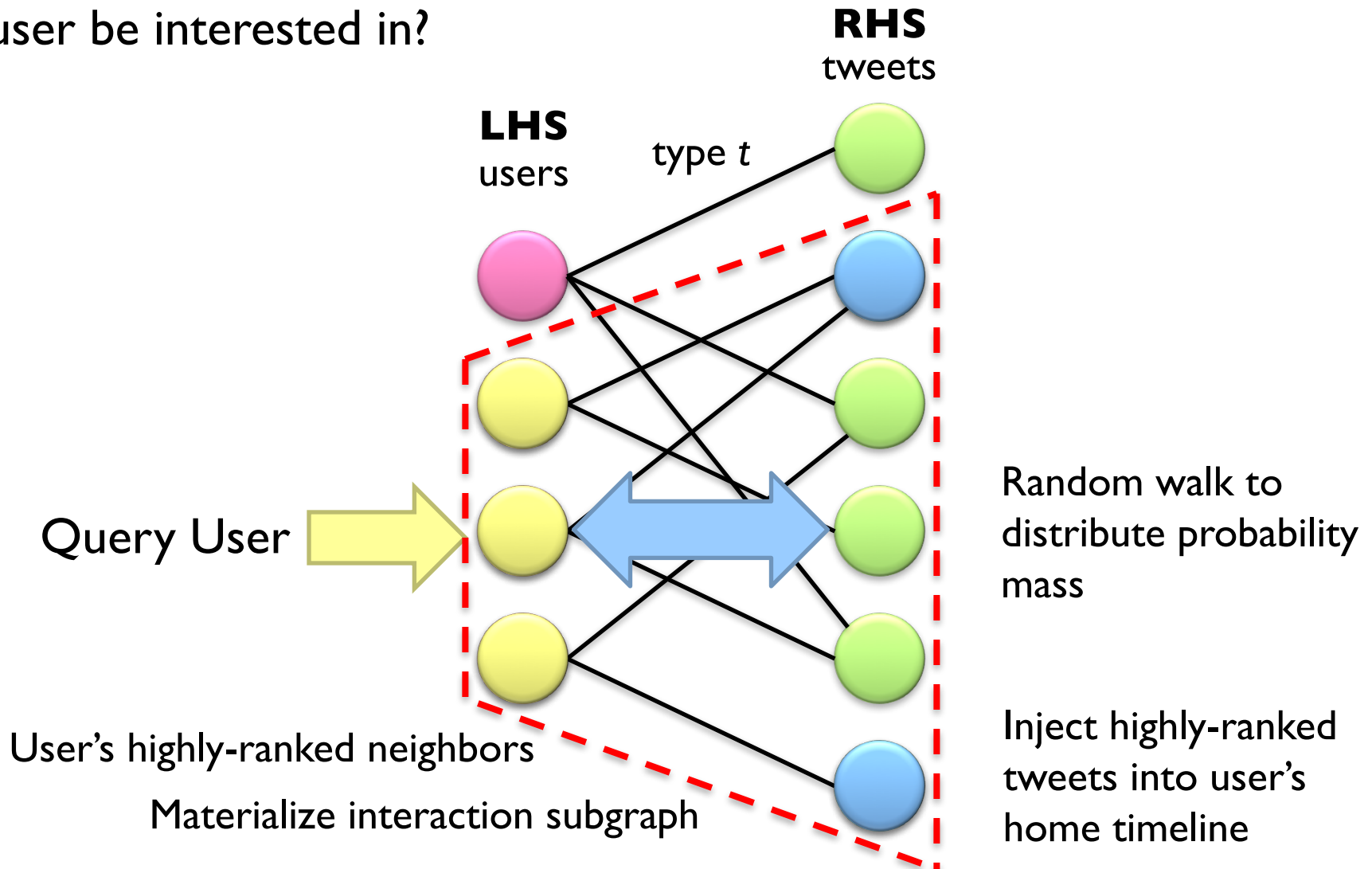
Marginally better results at the cost of space – not worthwhile

Note: design supports revisiting these choices

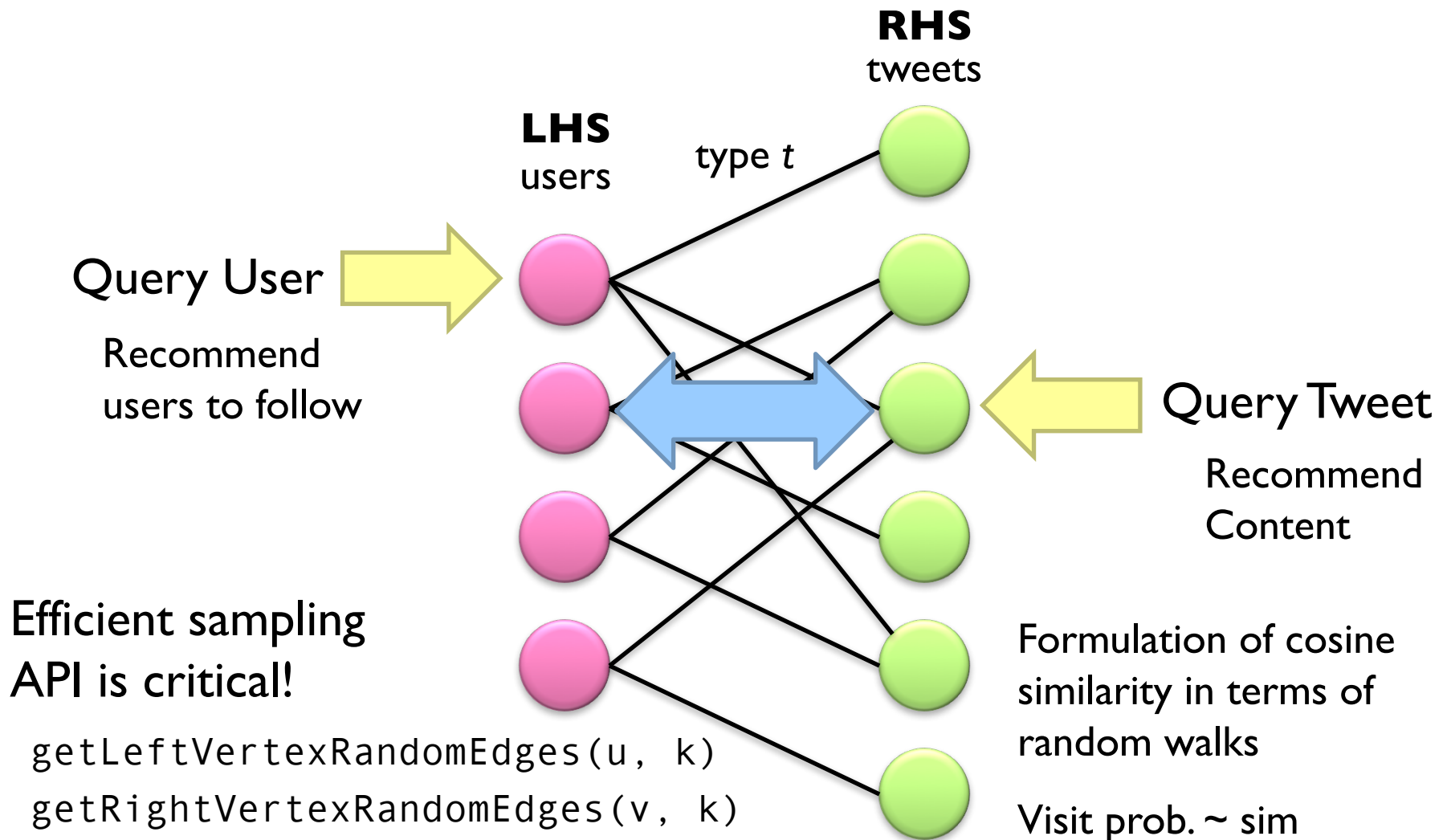


Recommendation Algorithm: Subgraph SALSA

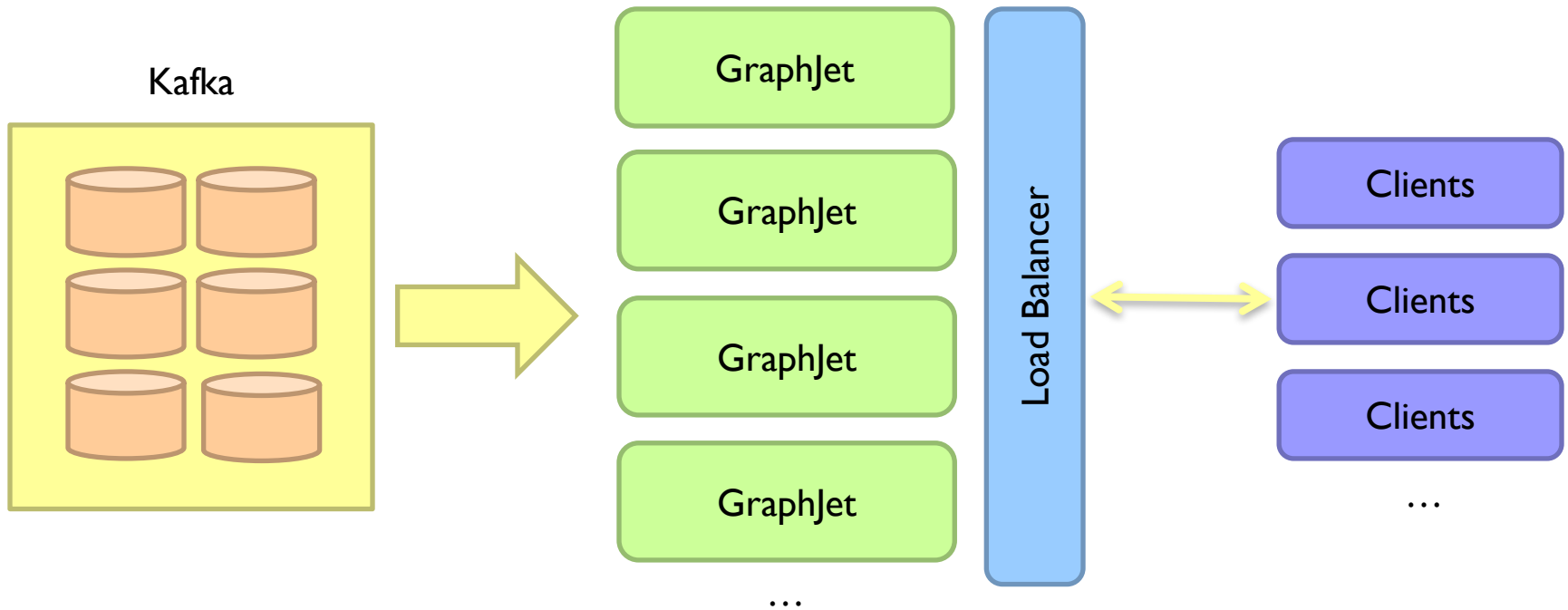
What tweets might a user be interested in?



Recommendation Algorithm: Similarity Query



Deployment Architecture



Production Status

Started serving production traffic early 2014

Dual Intel Xeon 6-cores (E5-2620 v2) at 2.1 GHz

Cold startup: ingestion at $O(10^6)$ edges per sec from Kafka

Steady state: ingestion at $O(10^4)$ edges per sec

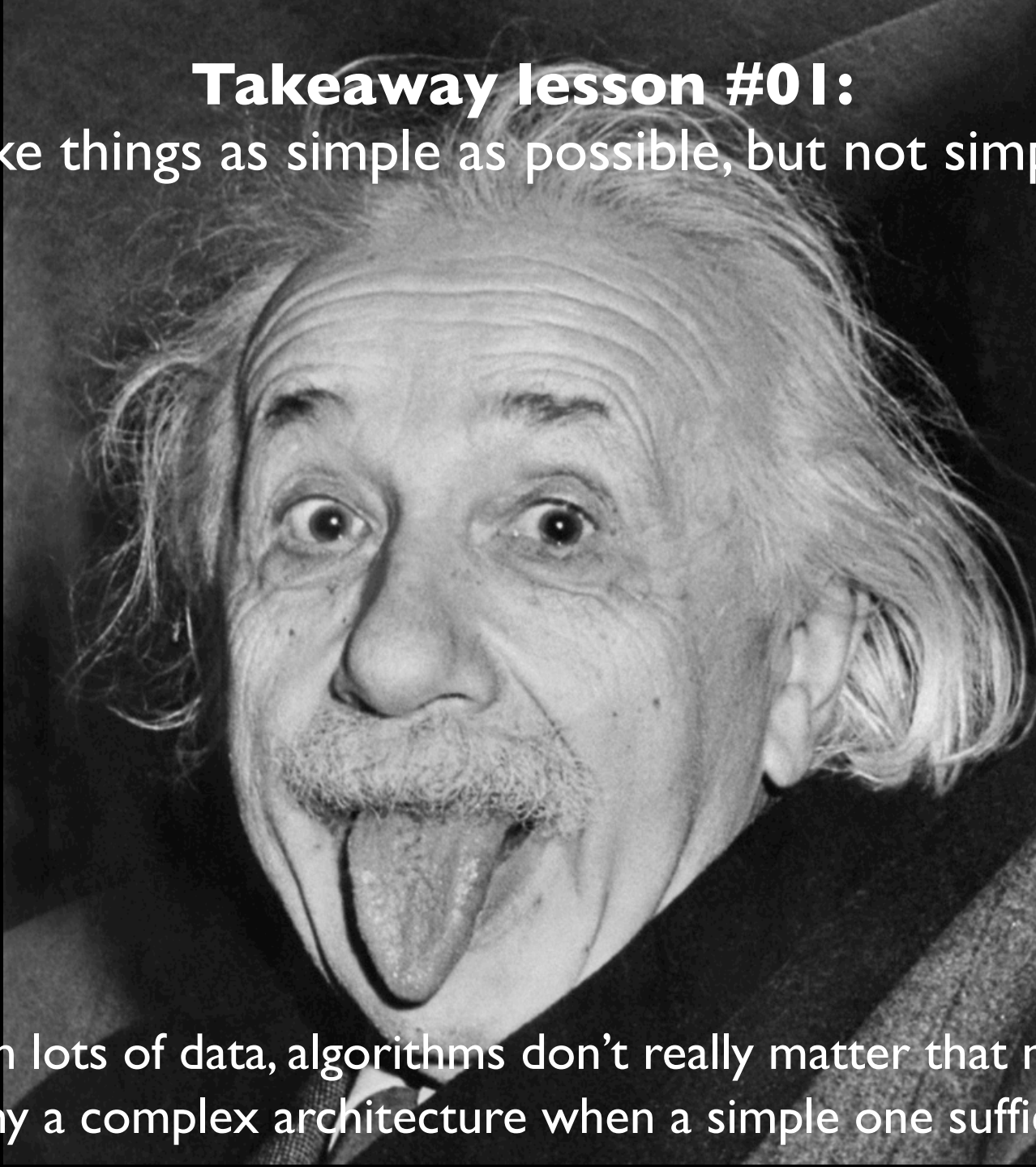
Space usage: $O(10^9)$ edges in < 30 GB

Sample recommendation algorithm: subgraph SALSA

500 QPS, p50 = 19ms, p99 = 33ms

Takeaway lesson #01:

Make things as simple as possible, but not simpler.



With lots of data, algorithms don't really matter that much
Why a complex architecture when a simple one suffices?



Takeaway lesson #10:
Constraints aren't always technical.

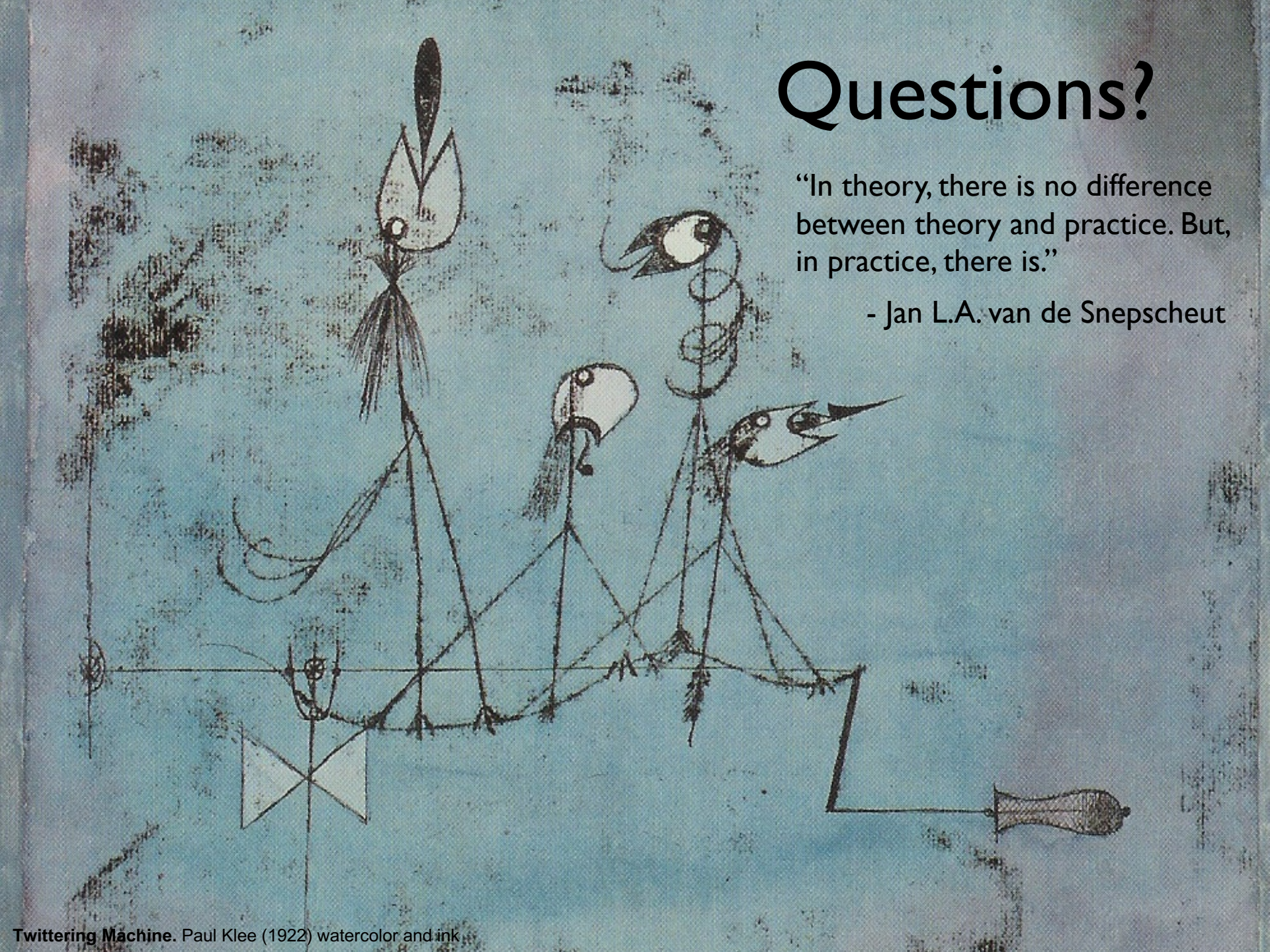


Takeaway lesson #11:
Visiting and revisiting design decisions

Questions?

“In theory, there is no difference between theory and practice. But, in practice, there is.”

- Jan L.A. van de Snepscheut



Twittering Machine. Paul Klee (1922) watercolor and ink