

Data-Intensive Distributed Computing

CS 451/651 (Fall 2018)

Part I: MapReduce Algorithm Design (3/4)
September 13, 2018

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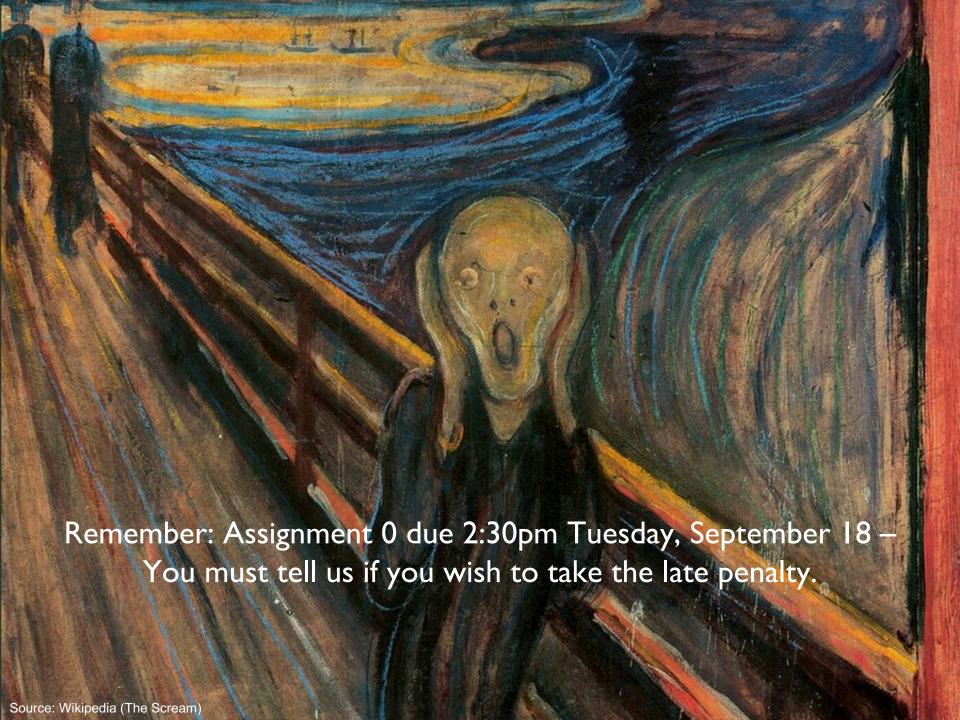
University of Waterloo

These slides are available at http://lintool.github.io/bigdata-2018f/

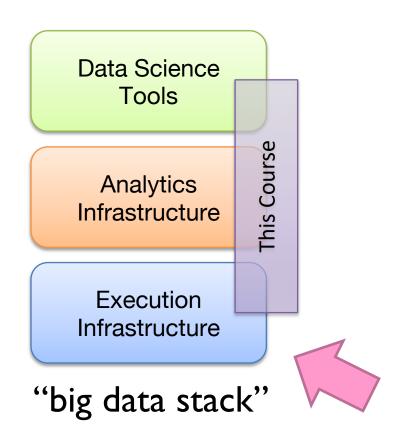


Agenda for Today

Cloud computing
Datacenter architectures
Hadoop cluster architecture
MapReduce physical execution



Today





The best thing since sliced bread?

Before clouds...

Grids

Connection machine

Vector supercomputers

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Cloud computing means many different things:

Big data
Rebranding of web 2.0
Utility computing
Everything as a service

Rebranding of web 2.0

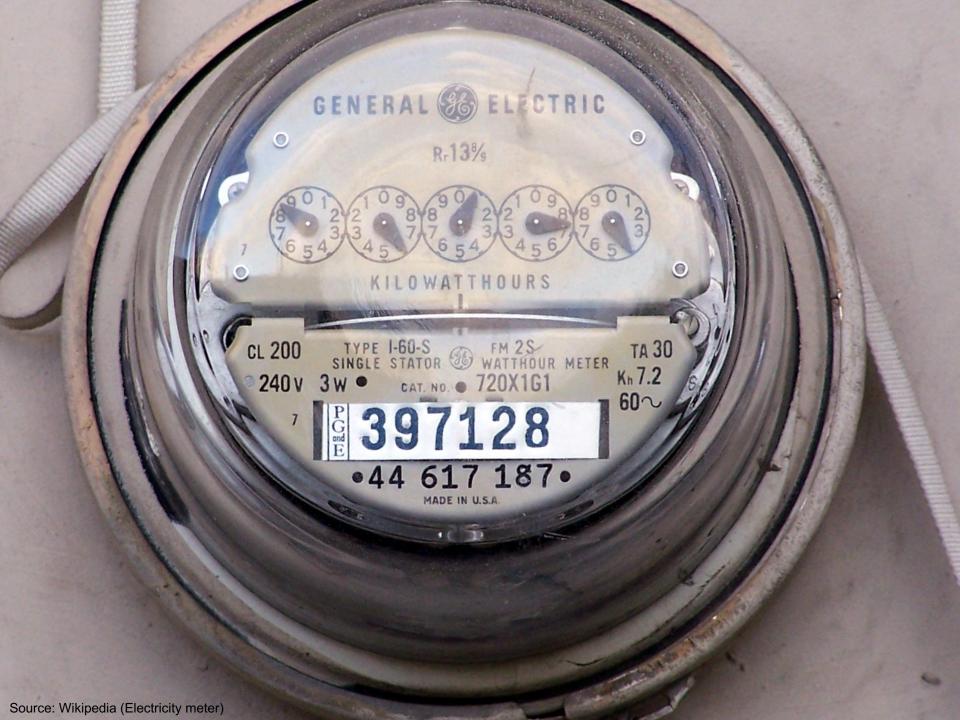
Rich, interactive web applications

Clouds refer to the servers that run them Javascript! (ugh)

Examples: Facebook, YouTube, Gmail, ...

"The network is the computer": take two

User data is stored "in the clouds"
Rise of the tablets, smartphones, etc. ("thin clients")
Browser is the OS



Utility Computing

What?

Computing resources as a metered service ("pay as you go")

Why?

Cost: capital vs. operating expenses

Scalability: "infinite" capacity

Elasticity: scale up or down on demand

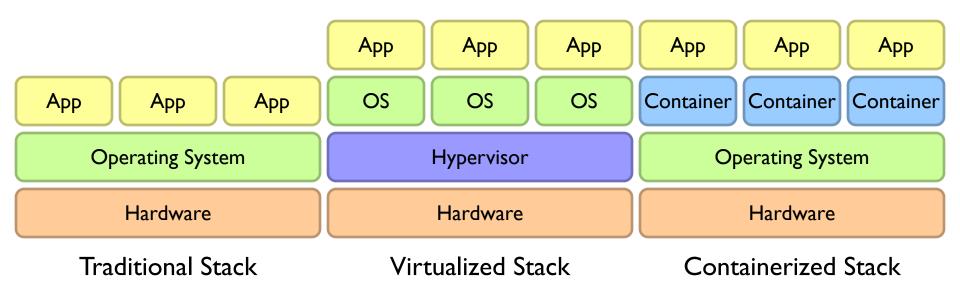
Does it make sense?

Benefits to cloud users
Business case for cloud providers

I think there is a world market for about five computers.



Evolution of the Stack



Everything as a Service

Infrastructure as a Service (laaS)

Why buy machines when you can rent them instead? Examples: Amazon EC2, Microsoft Azure, Google Compute

Platform as a Service (PaaS)

Give me a nice platform and take care of maintenance, upgrades, ... Example: Google App Engine

Software as a Service (SaaS)

Just run the application for me! Example: Gmail, Salesforce

Everything as a Service

Database as a Service

Run a database for me

Examples: Amazon RDS, Microsoft Azure SQL

Search as a Service

Run a search engine for me

Example: Amazon Elasticsearch Service

Function as a Service

Run this function for me

Example: Amazon Lambda, Google Cloud Functions

Who cares?

A source of problems...

Cloud-based services generate big data
Clouds make it easier to start companies that generate big data

As well as a solution...

Ability to provision clusters on-demand in the cloud Commoditization and democratization of big data capabilities



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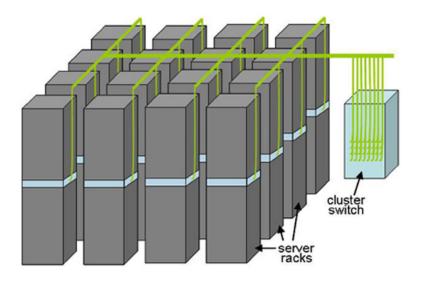




Building Blocks





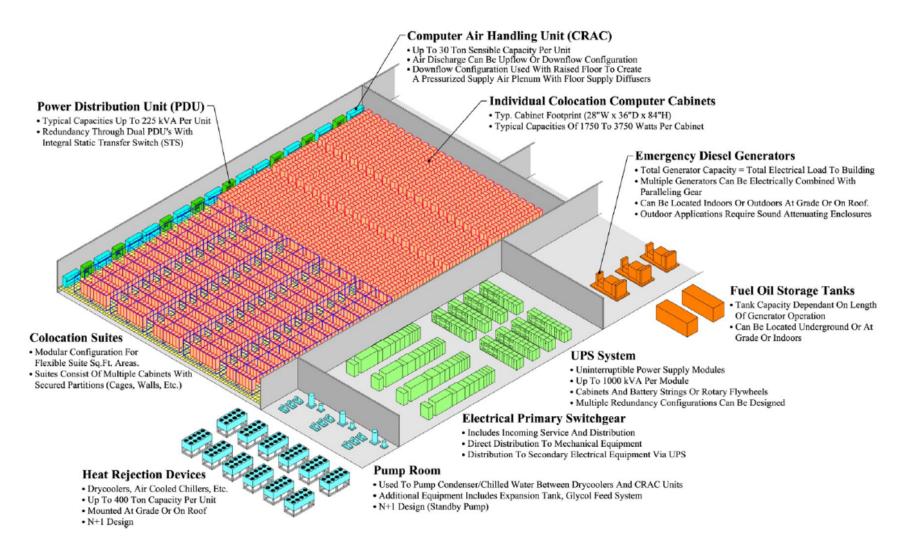






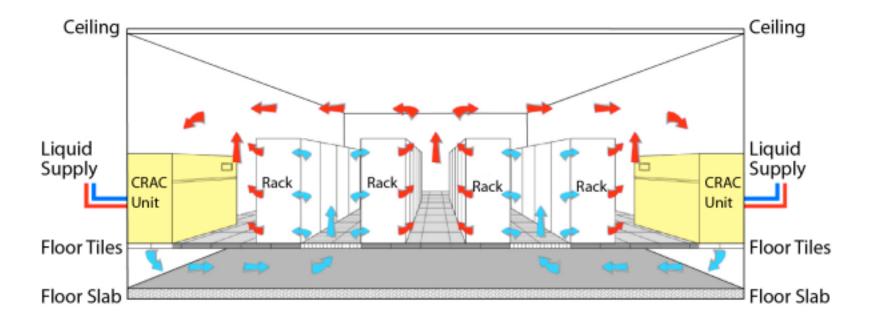


Anatomy of a Datacenter



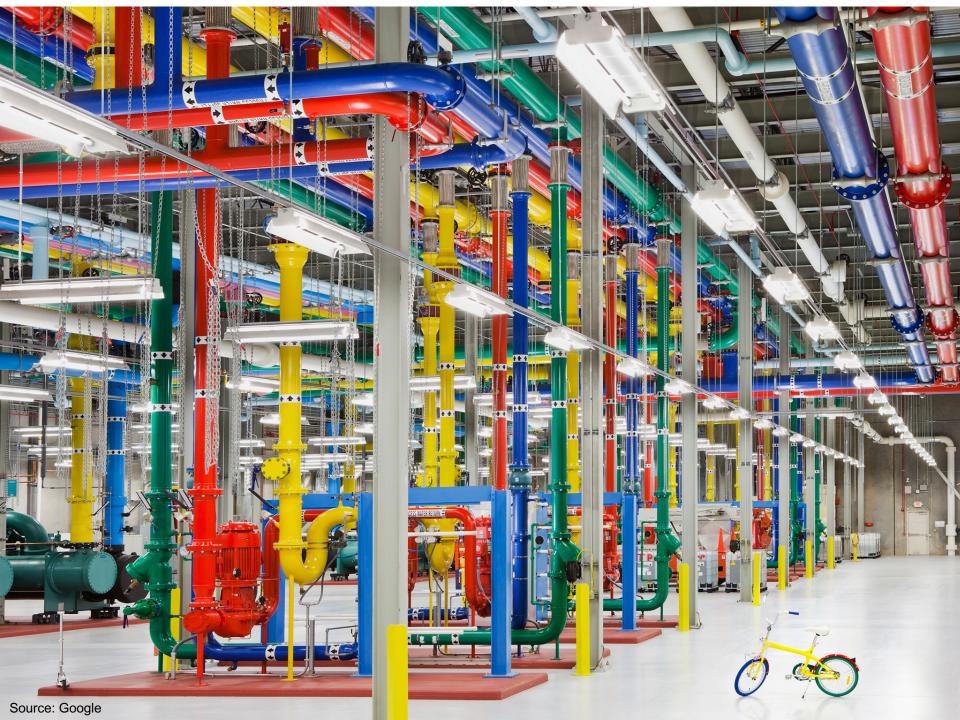
Datacenter cooling

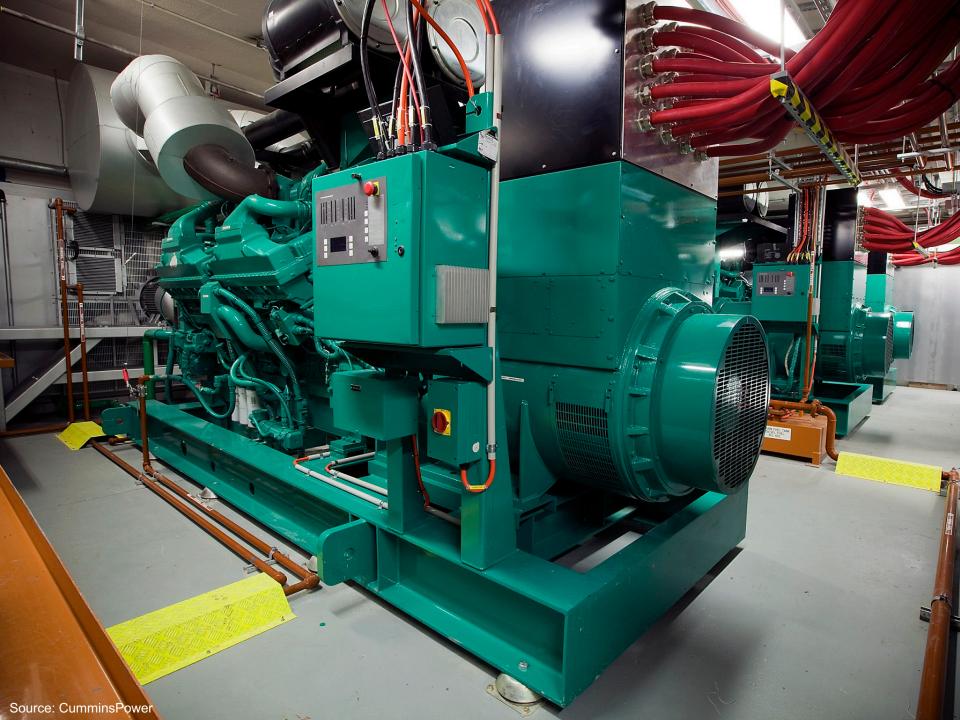
What's a computer?



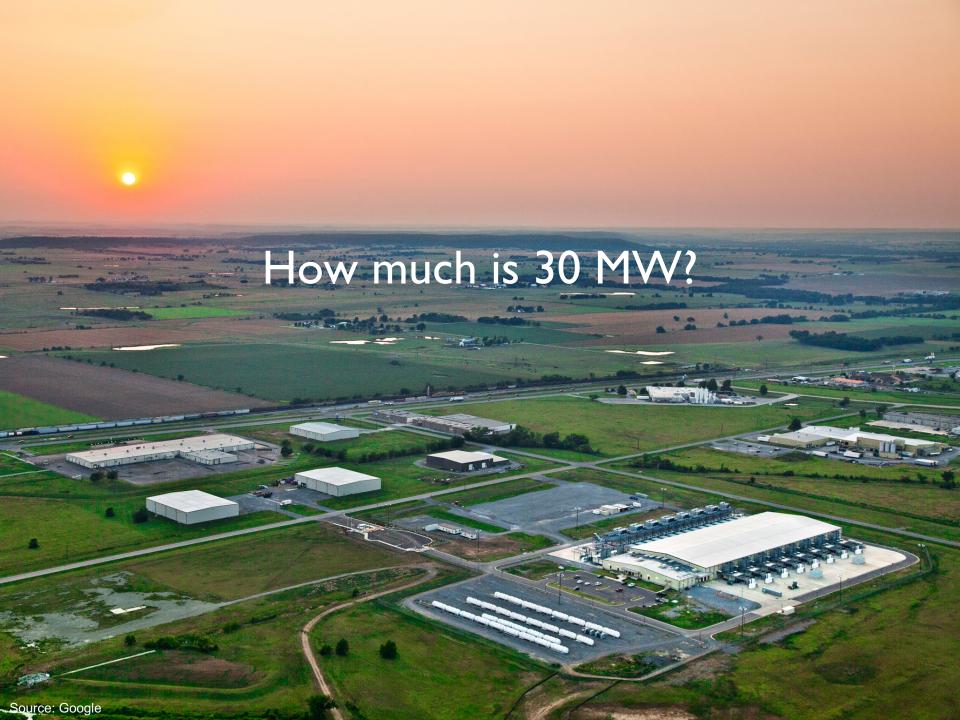
Source: Barroso and Urs Hölzle (2013)



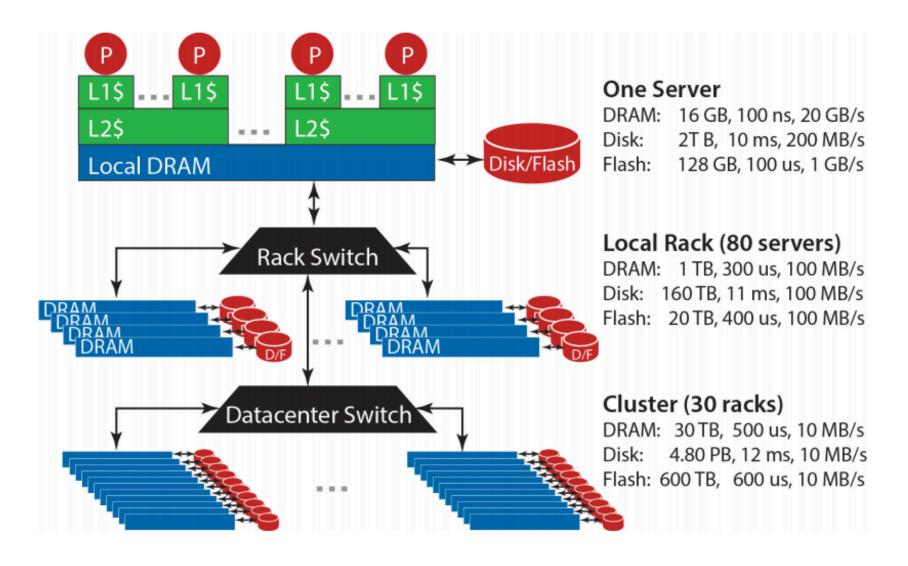








Datacenter Organization



Source: Barroso and Urs Hölzle (2013)

The datacenter is the computer!

It's all about the right level of abstraction

Moving beyond the von Neumann architecture What's the "instruction set" of the datacenter computer?

Hide system-level details from the developers

No more race conditions, lock contention, etc. No need to explicitly worry about reliability, fault tolerance, etc.

Separating the what from the how

Developer specifies the computation that needs to be performed Execution framework ("runtime") handles actual execution

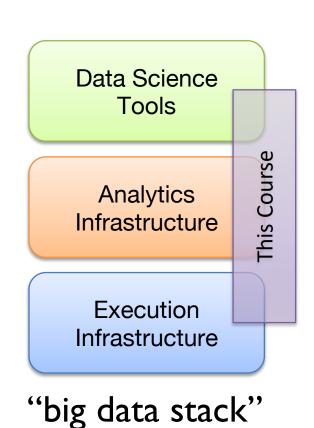


Mechanical Sympathy

"You don't have to be an engineer to be be a racing driver, but you do have to have mechanical sympathy"

- Formula One driver Jackie Stewart





Intuitions of time and space

How long does it take to read 100 TBs from 100 hard drives? Now, what about SSDs?

How long will it take to exchange 1b key-value pairs:

Between machines on the same rack?

Between datacenters across the Atlantic?

Storage Hierarchy

Remote Machine

Different Datacenter

Remote Machine Different Rack

Remote Machine Same Rack

Local Machine
LI/L2/L3 cache, memory, SSD, magnetic disks
capacity, latency, bandwidth

Numbers Everyone Should Know According to Jeff Dean

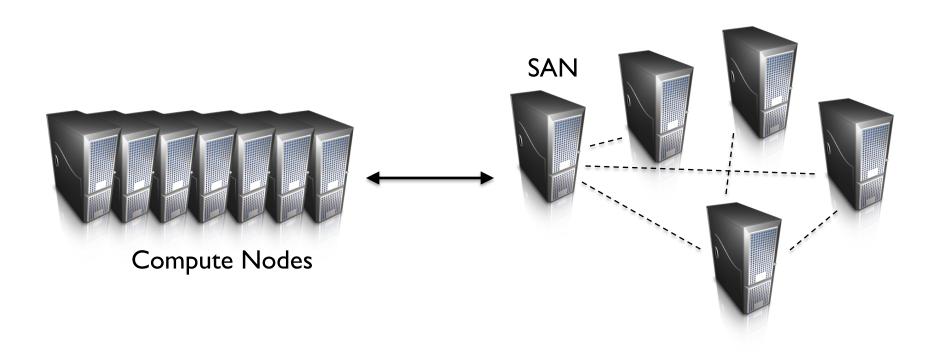
| L1 cache reference | 0 . | .5 ns |
|-------------------------------------|-------------|-------|
| Branch mispredict | 5 | ns |
| L2 cache reference | 7 | ns |
| Mutex lock/unlock | 100 | ns |
| Main memory reference | 100 | ns |
| Compress 1K bytes with Zippy | 10,000 | ns |
| Send 2K bytes over 1 Gbps network | 20,000 | ns |
| Read 1 MB sequentially from memory | 250,000 | ns |
| Round trip within same datacenter | 500,000 | ns |
| Disk seek | 10,000,000 | ns |
| Read 1 MB sequentially from network | 10,000,000 | ns |
| Read 1 MB sequentially from disk | 30,000,000 | ns |
| Send packet CA->Netherlands->CA | 150,000,000 | ns |



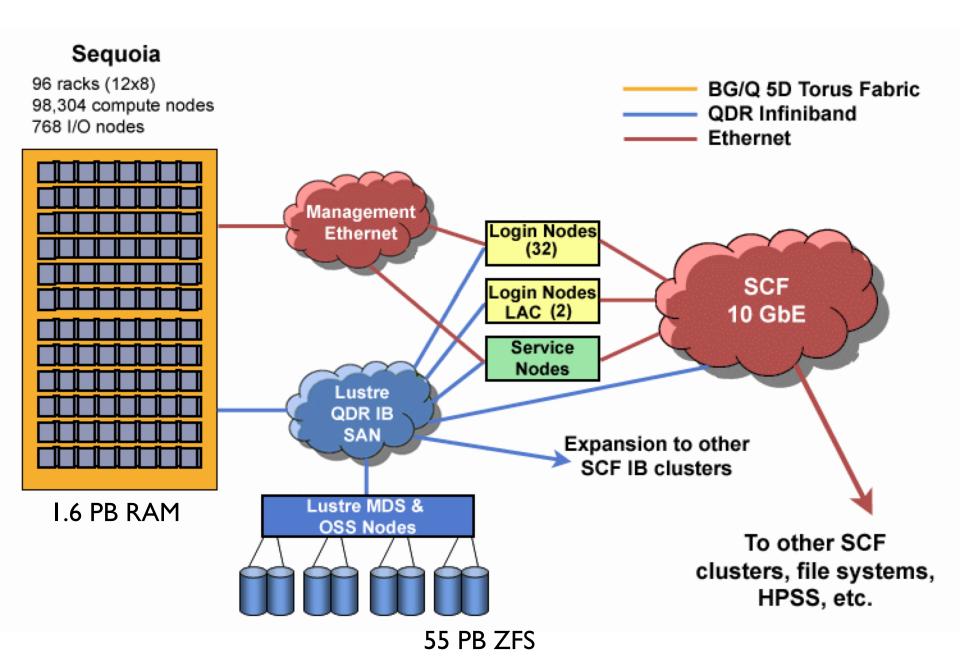


How do we get data to the workers?

Let's consider a typical supercomputer...

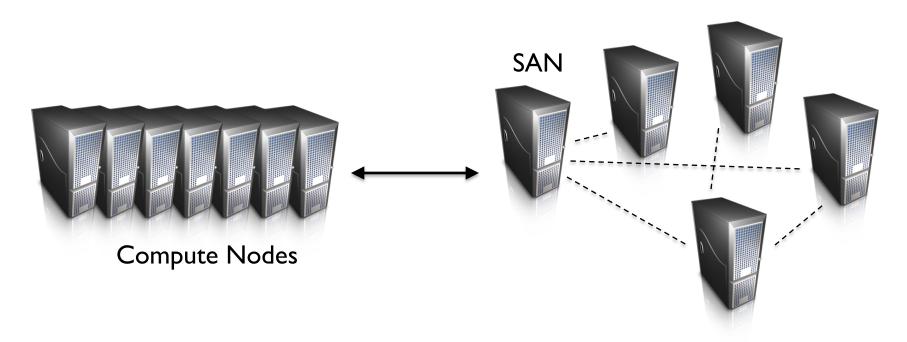






Source: LLNL

Compute-Intensive vs. Data-Intensive



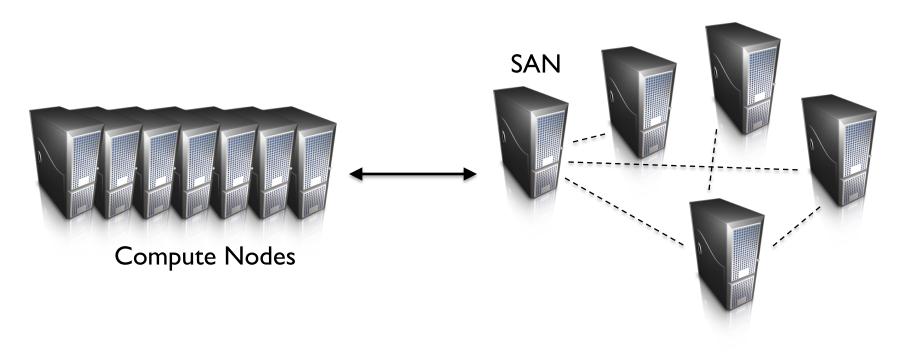
Why does this make sense for compute-intensive tasks? What's the issue for data-intensive tasks?

What's the solution?

Don't move data to workers... move workers to the data!

Key idea: co-locate storage and compute

Start up worker on nodes that hold the data



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Don't move data to workers... move workers to the data!

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We need a distributed file system for managing this

GFS (Google File System) for Google's MapReduce HDFS (Hadoop Distributed File System) for Hadoop

GFS: Assumptions

Commodity hardware over "exotic" hardware Scale "out", not "up"

High component failure rates

Inexpensive commodity components fail all the time

"Modest" number of huge files Multi-gigabyte files are common, if not encouraged

Files are write-once, mostly appended to Logs are a common case

Large streaming reads over random access

Design for high sustained throughput over low latency

GFS: Design Decisions

Files stored as chunks

Fixed size (64MB)

Reliability through replication

Each chunk replicated across 3+ chunkservers

Single master to coordinate access and hold metadata

Simple centralized management

No data caching

Little benefit for streaming reads over large datasets

Simplify the API: not POSIX!

Push many issues onto the client (e.g., data layout)

HDFS = GFS clone (same basic ideas)

From GFS to HDFS

Terminology differences:

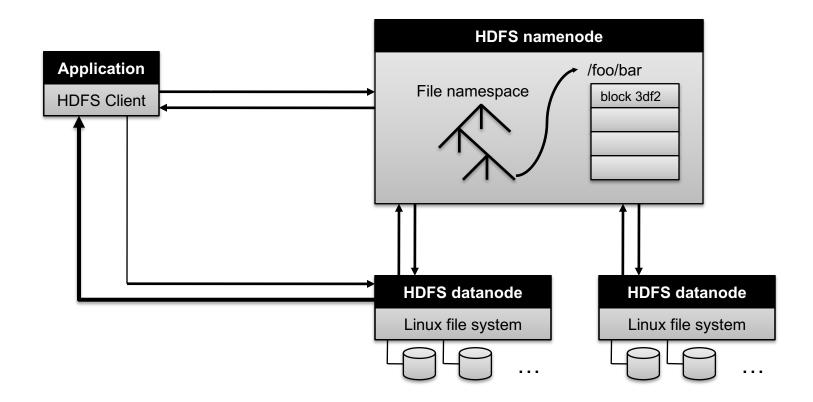
GFS master = Hadoop namenode GFS chunkservers = Hadoop datanodes

Implementation differences:

Different consistency model for file appends
Implementation language
Performance

For the most part, we'll use Hadoop terminology...

HDFS Architecture



Namenode Responsibilities

Managing the file system namespace

Holds file/directory structure, file-to-block mapping, metadata (ownership, access permissions, etc.)

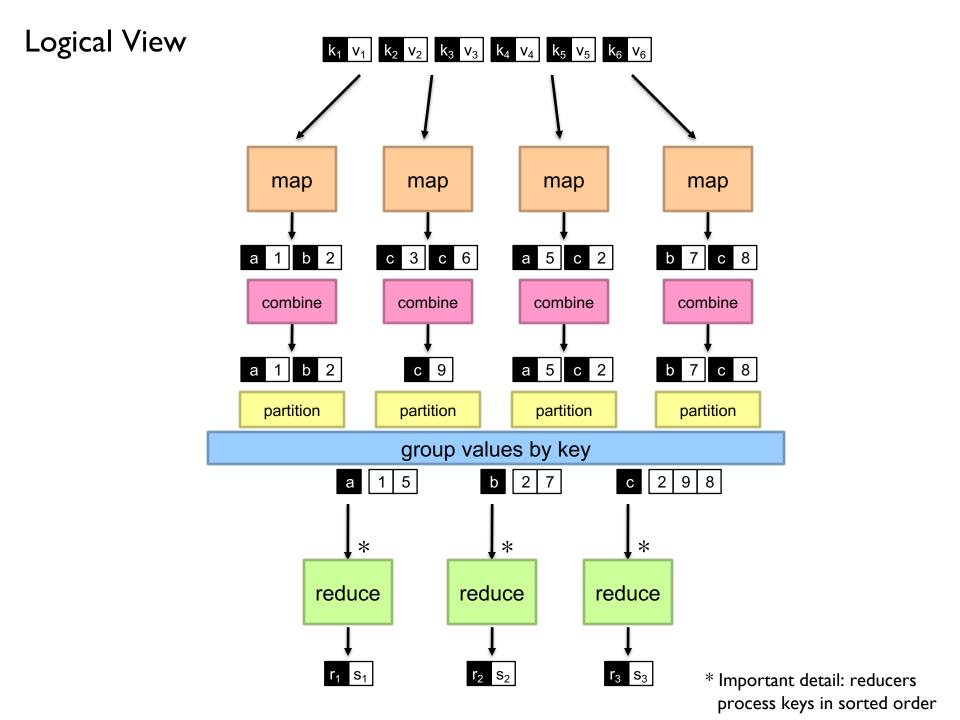
Coordinating file operations

Directs clients to datanodes for reads and writes

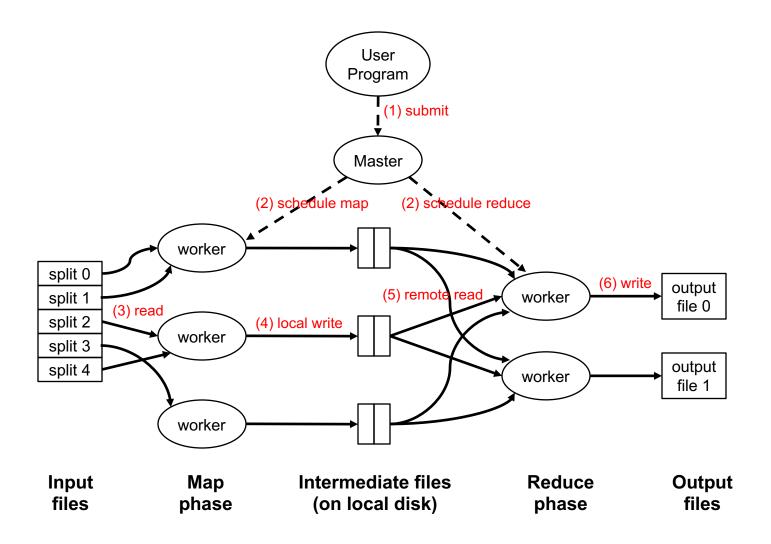
No data is moved through the namenode

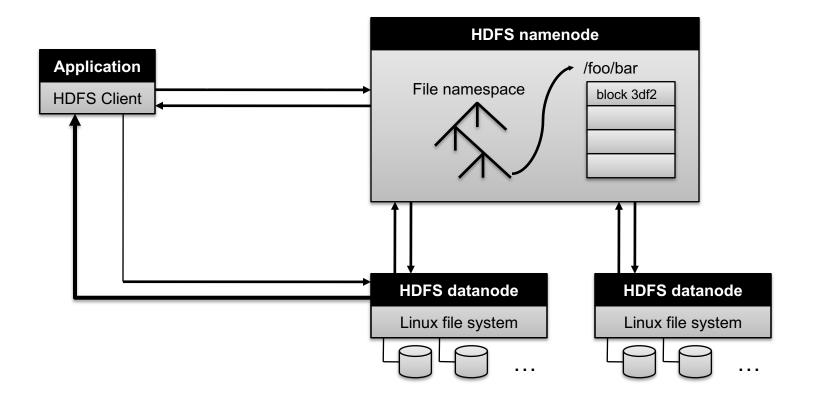
Maintaining overall health

Periodic communication with the datanodes
Block re-replication and rebalancing
Garbage collection

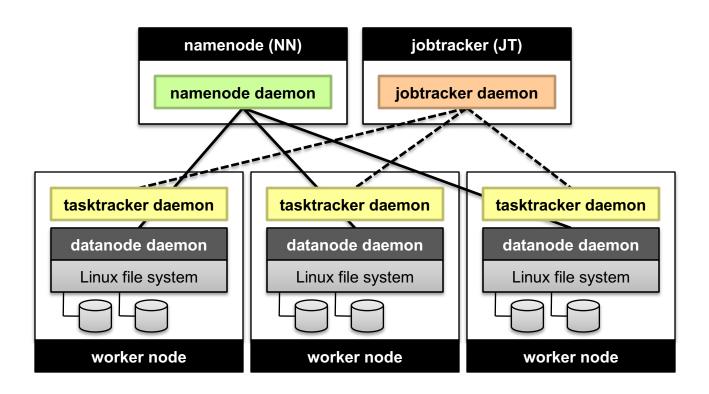


Physical View





Putting everything together...



Basic Cluster Components*

Namenode (NN)

Master for HDFS

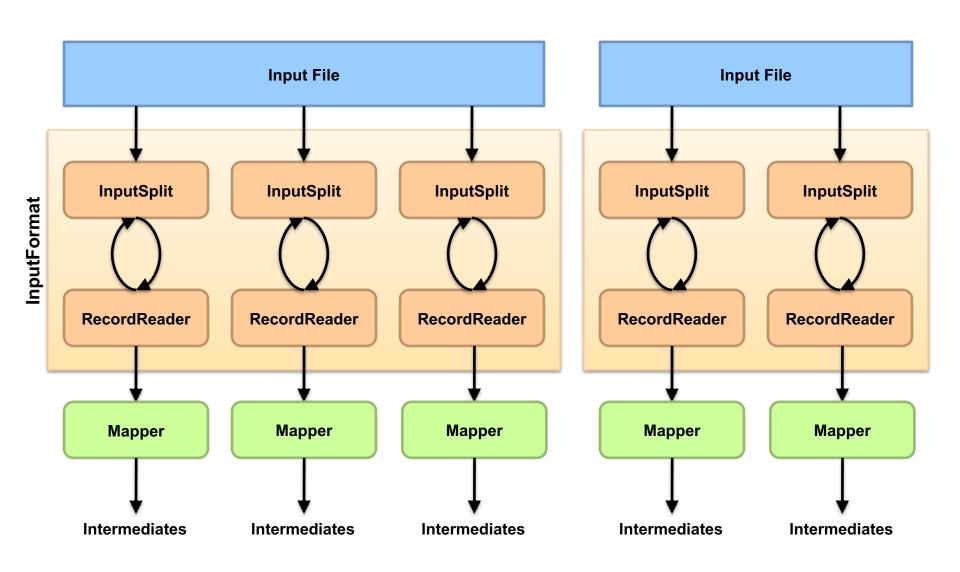
Jobtracker (JT)

Coordinator for MapReduce jobs

On each of the worker machines:

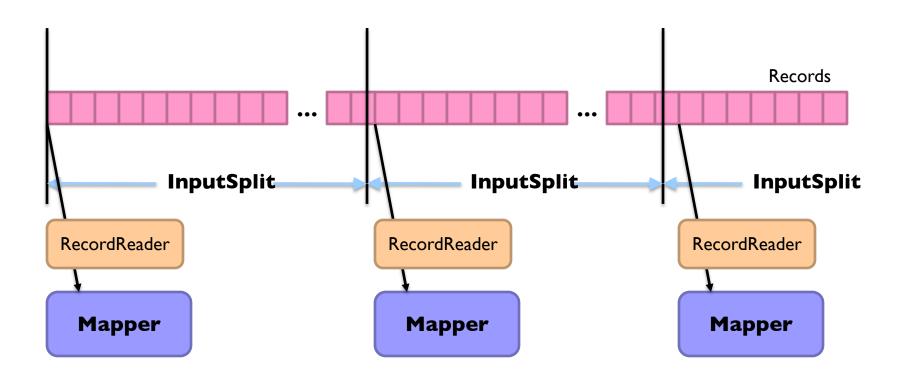
Tasktracker (TT): contains multiple task slots

Datanode (DN): serves HDFS data blocks

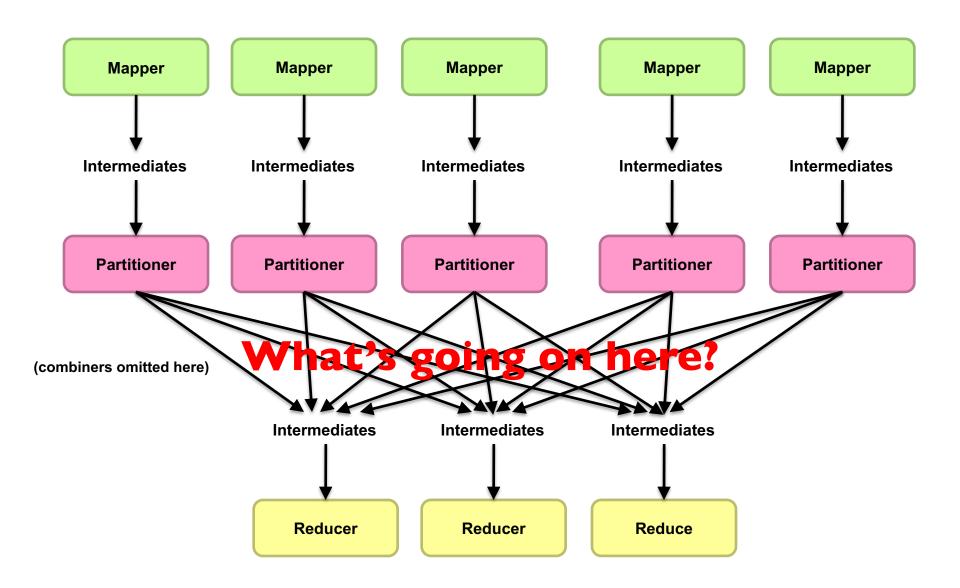


What are these input split?

Client



What are these input split?



Distributed Group By in MapReduce

Map side

Map outputs are buffered in memory in a circular buffer When buffer reaches threshold, contents are "spilled" to disk Spills are merged into a single, partitioned file (sorted within each partition) Combiner runs during the merges

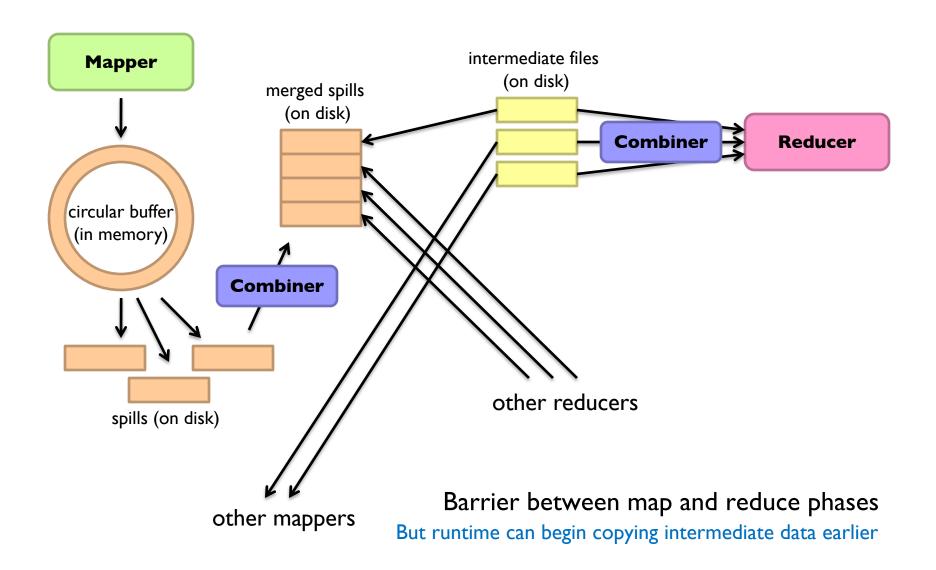
Reduce side

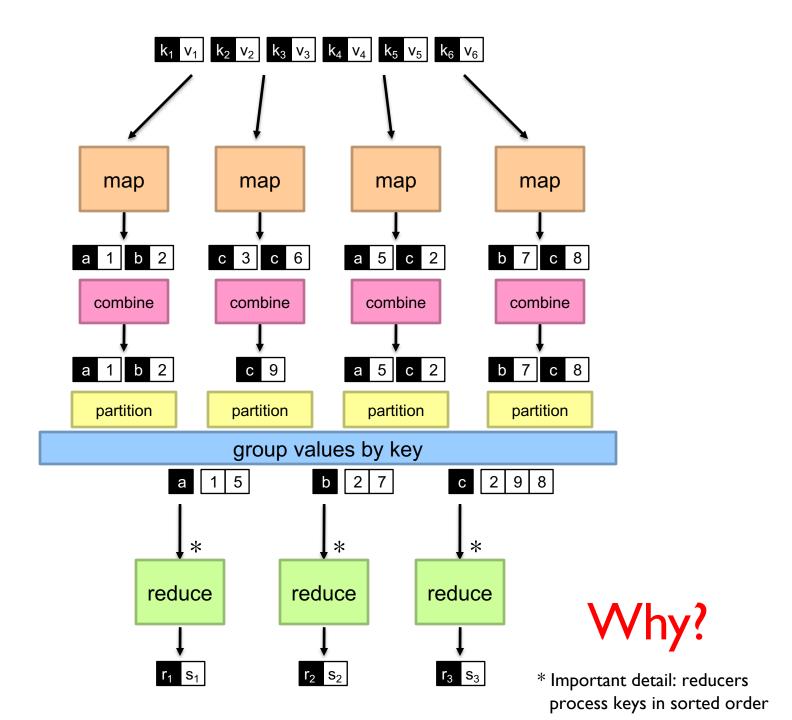
First, map outputs are copied over to reducer machine "Sort" is a multi-pass merge of map outputs (happens in memory and on disk)

Combiner runs during the merges

Final merge pass goes directly into reducer

Distributed Group By in MapReduce





Law of Leaky Abstractions

All non-trivial abstractions, to some degree, are leaky.

Joel Spolsky

