

Big Data Infrastructure

Session 10: Beyond MapReduce — Graph Processing

Jimmy Lin
University of Maryland
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Today's Agenda

- What makes graph processing hard?
- Graph processing frameworks
- Twitter case study

What makes graph processing hard?

- Lessons learned so far:
 - Partition
 - Replicate
 - Reduce cross-partition communication
- What makes MapReduce “work”?

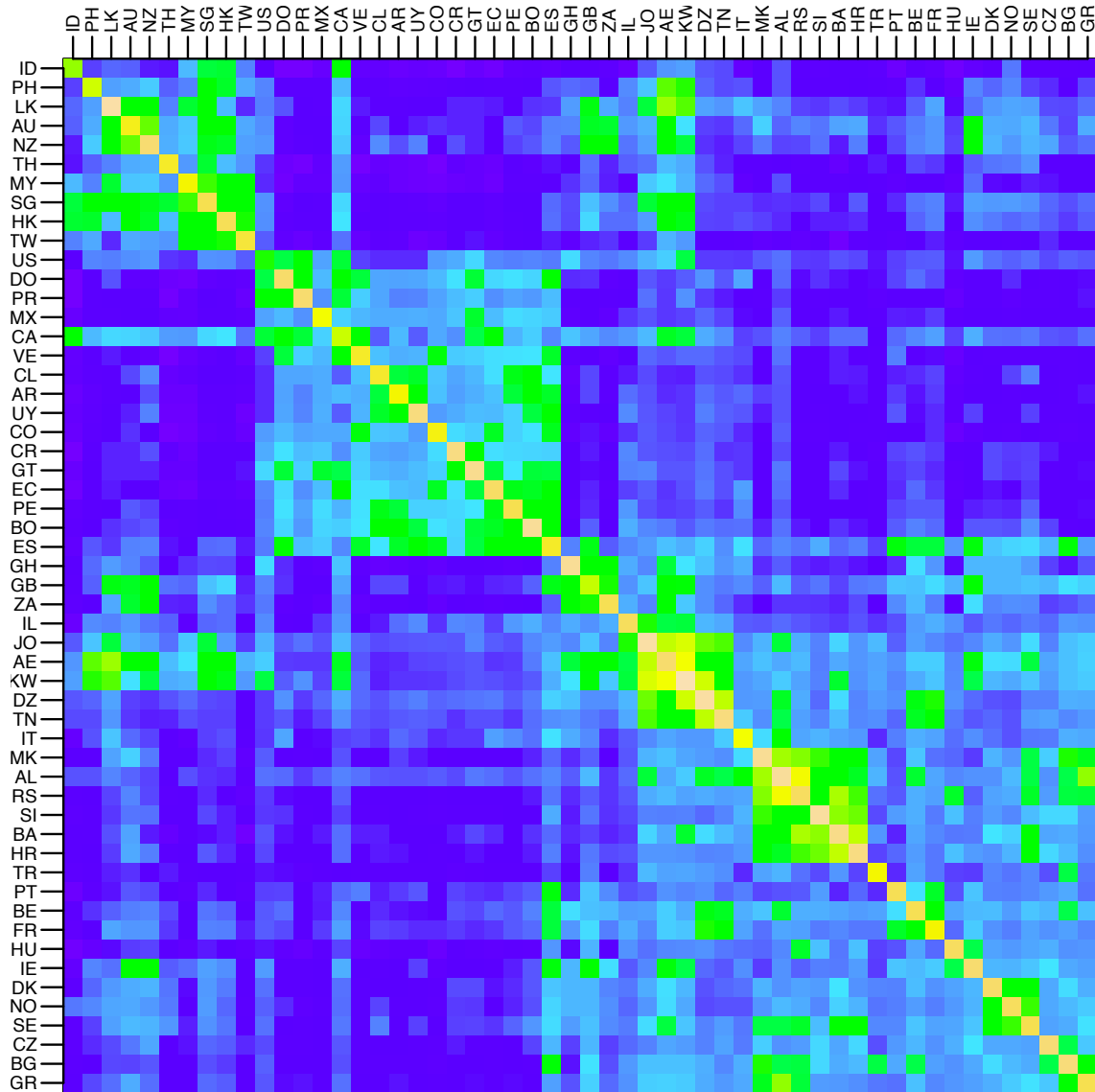
Characteristics of Graph Algorithms

- What are some common features of graph algorithms?
 - Graph traversals
 - Computations involving vertices and their neighbors
 - Passing information along graph edges
- What's the obvious idea?
 - Keep “neighborhoods” together!

Simple Partitioning Techniques

- Hash partitioning
- Range partitioning on some underlying linearization
 - Web pages: lexicographic sort of domain-reversed URLs
 - Social networks: sort by demographic characteristics

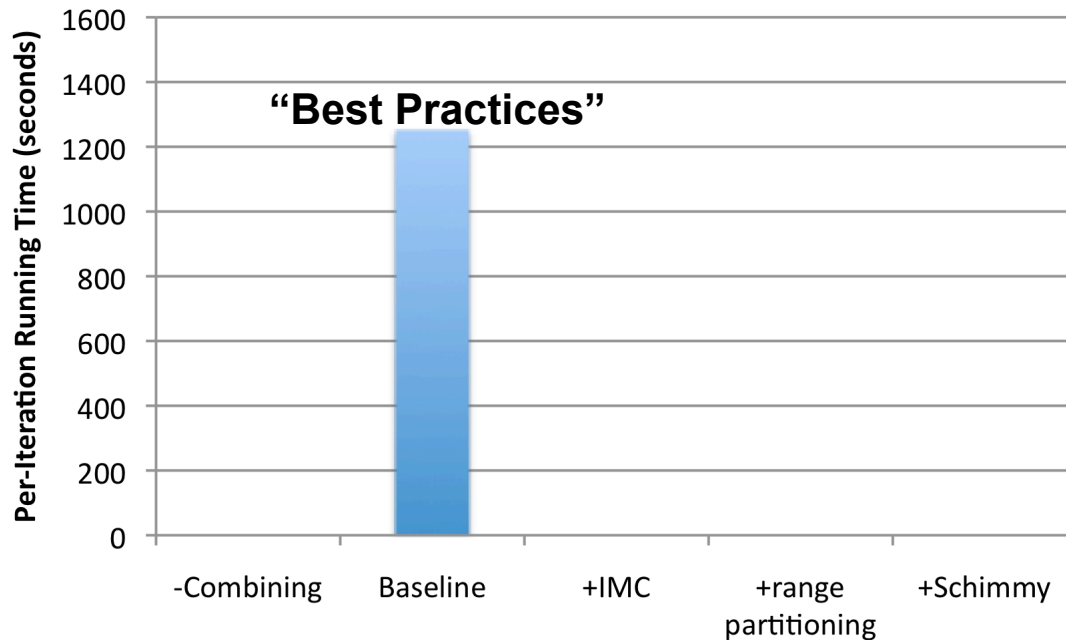
Country Structure in Facebook



Analysis of 721 million active users (May 2011)

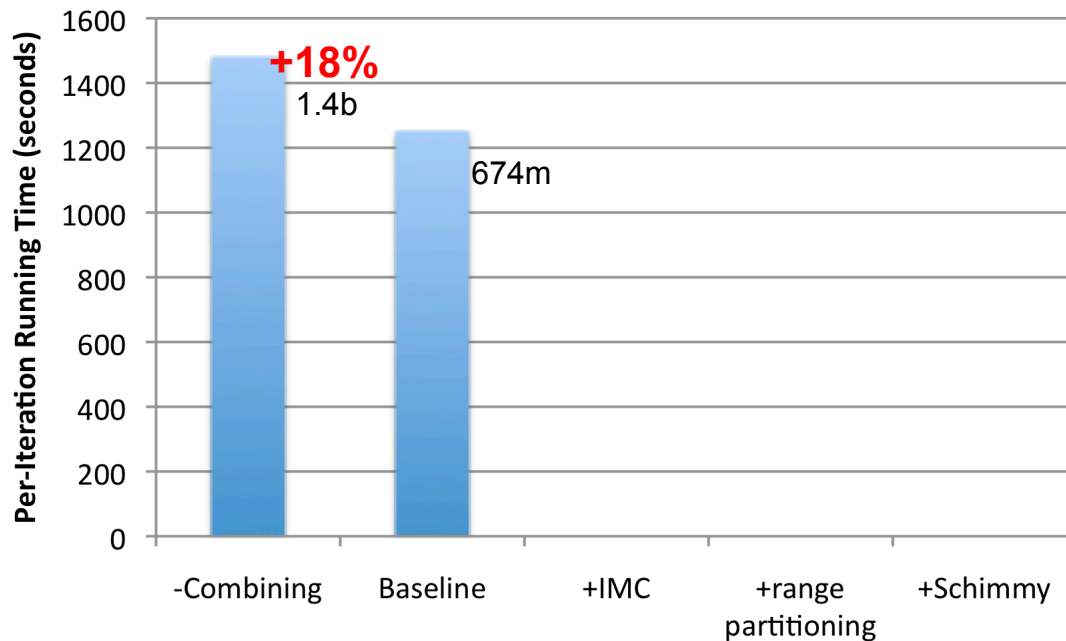
54 countries w/ >1m active users, >50% penetration

How much difference does it make?



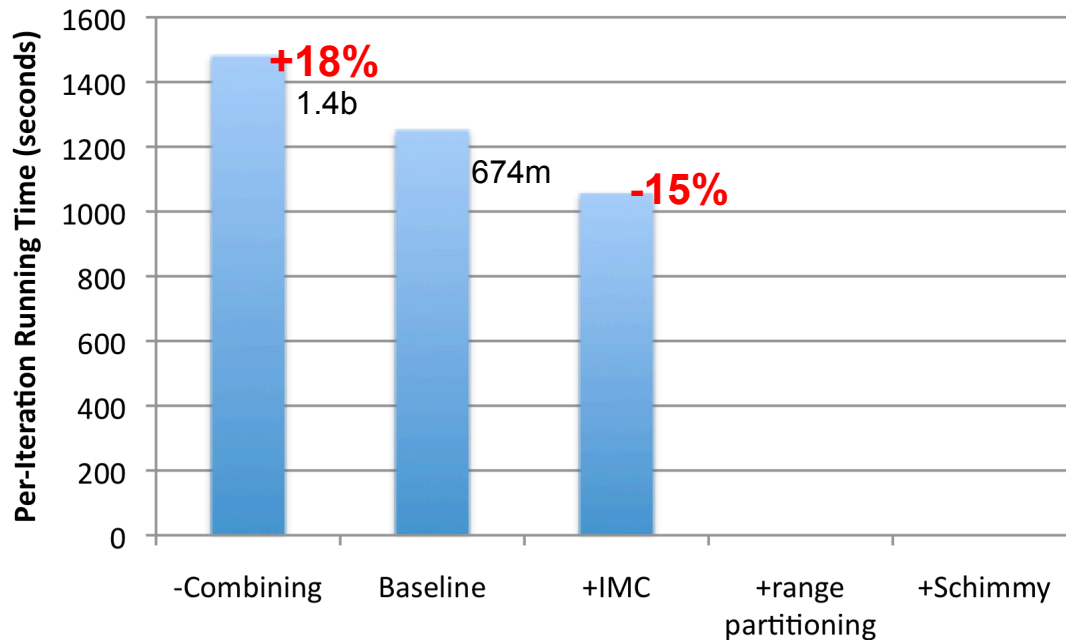
PageRank over webgraph
(40m vertices, 1.4b edges)

How much difference does it make?



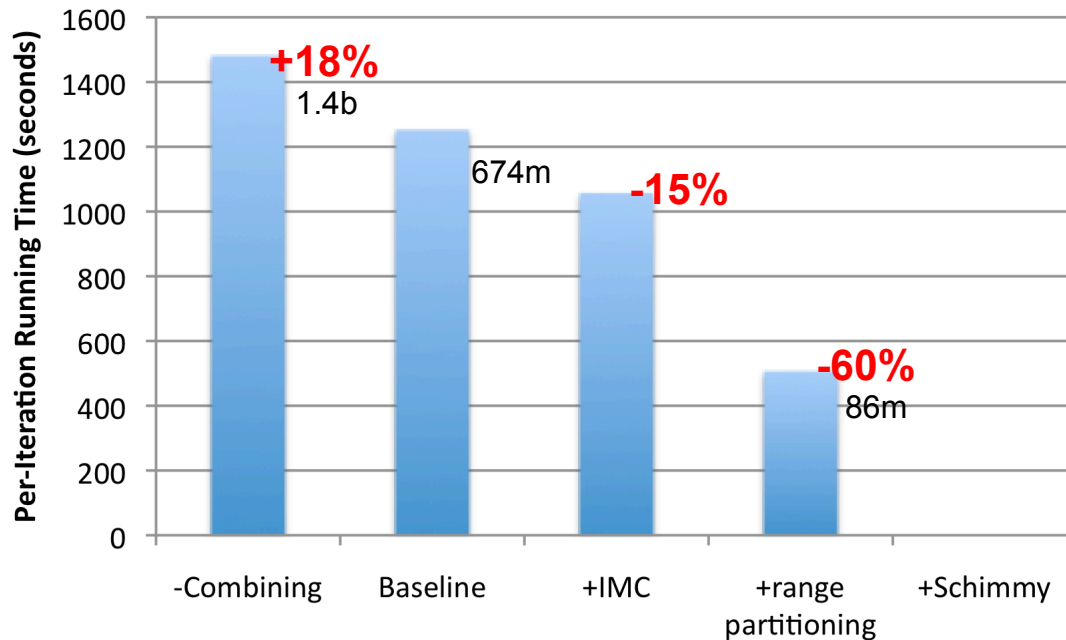
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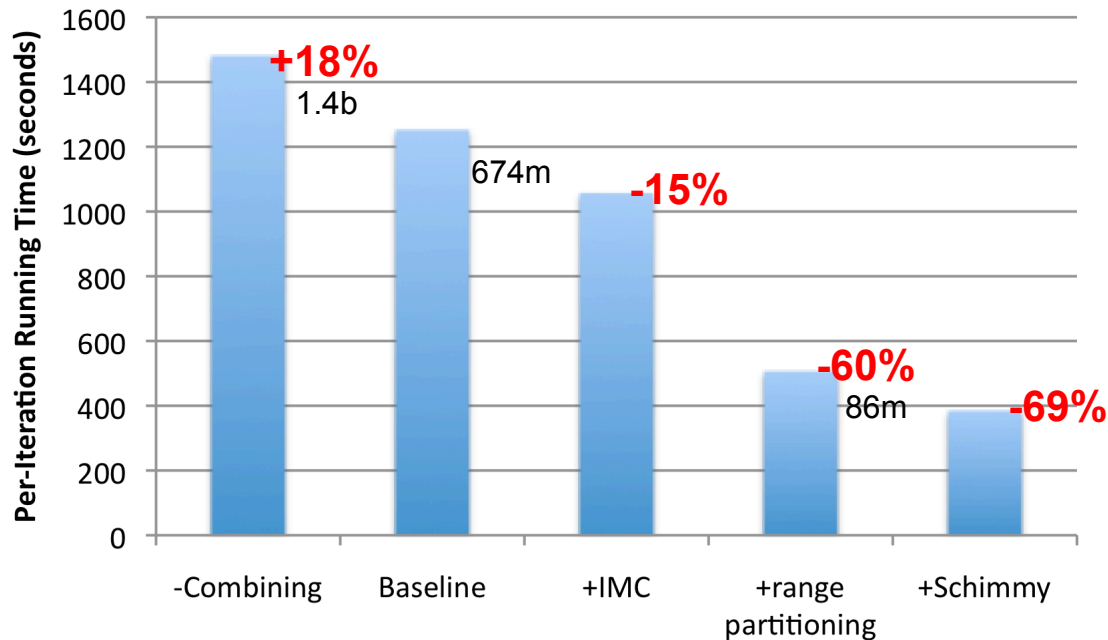
PageRank over webgraph
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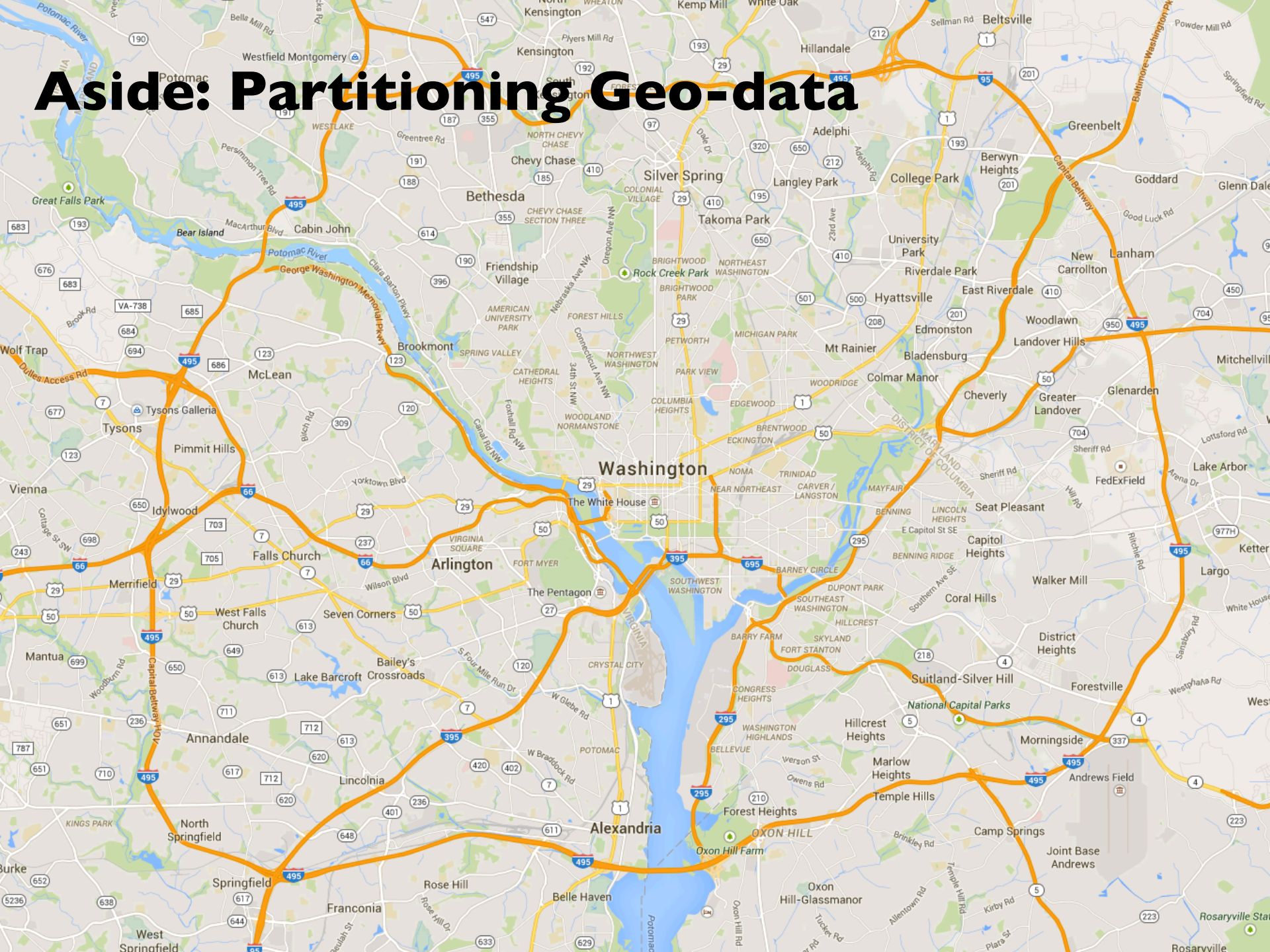
PageRank over webgraph
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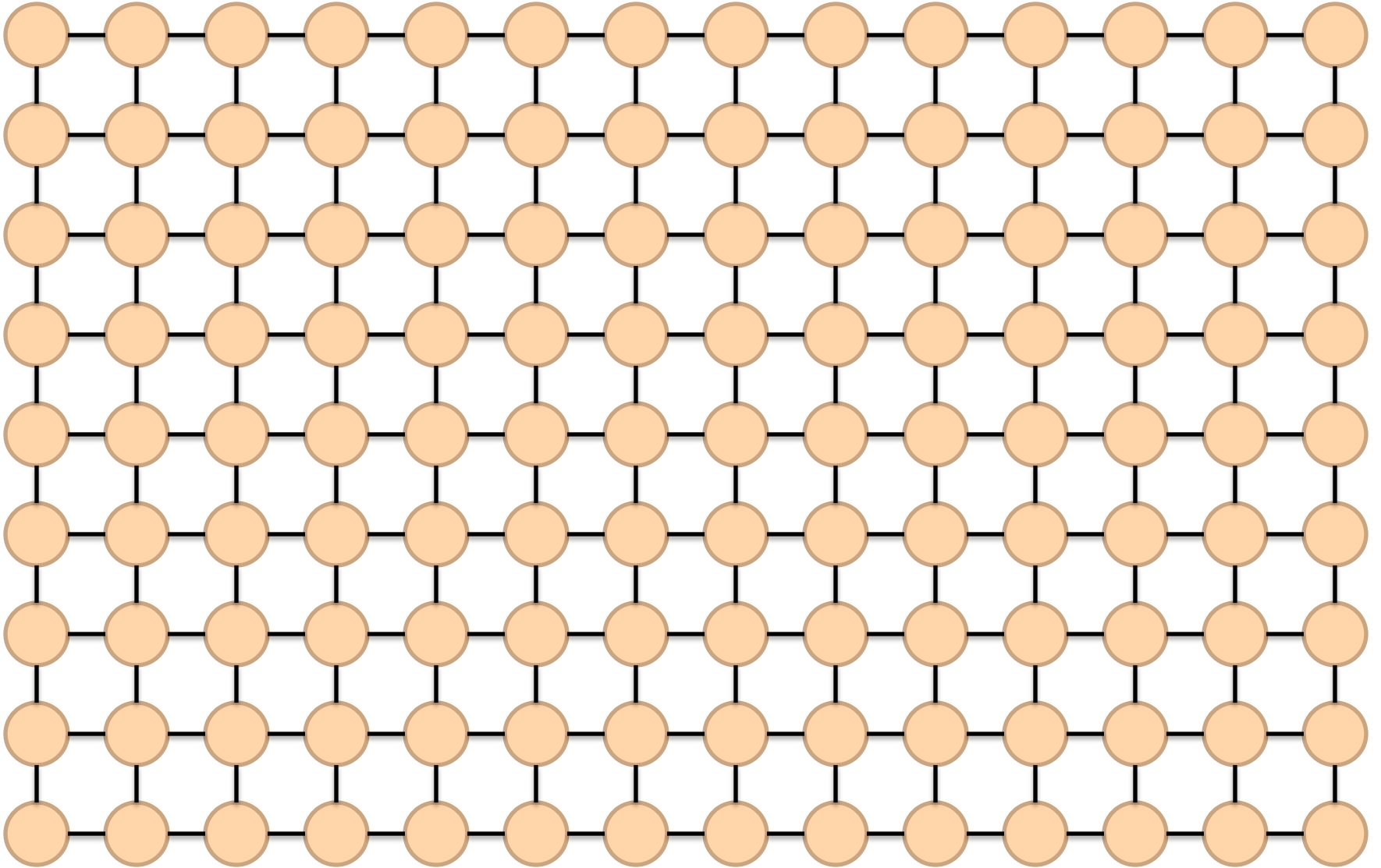


PageRank over webgraph
(40m vertices, 1.4b edges)

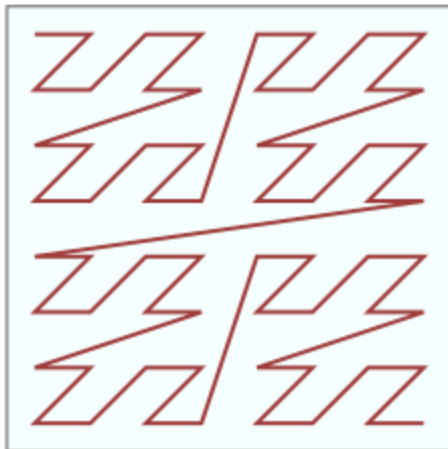
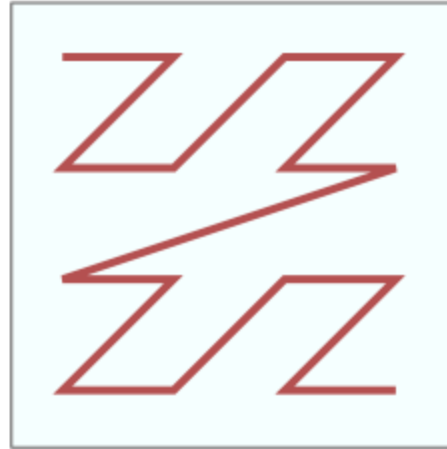
Aside: Partitioning Geo-data



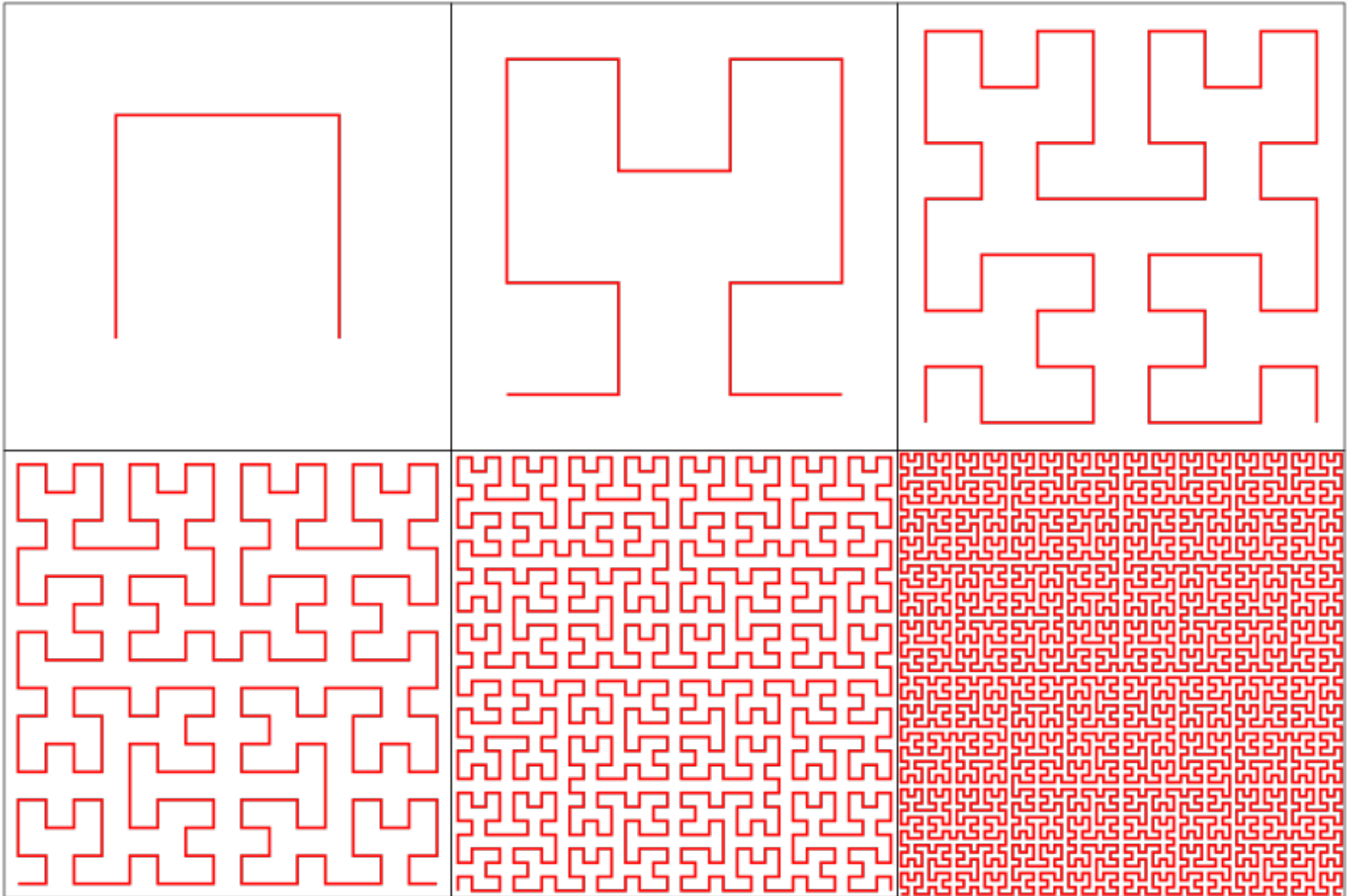
Geo-data = regular graph



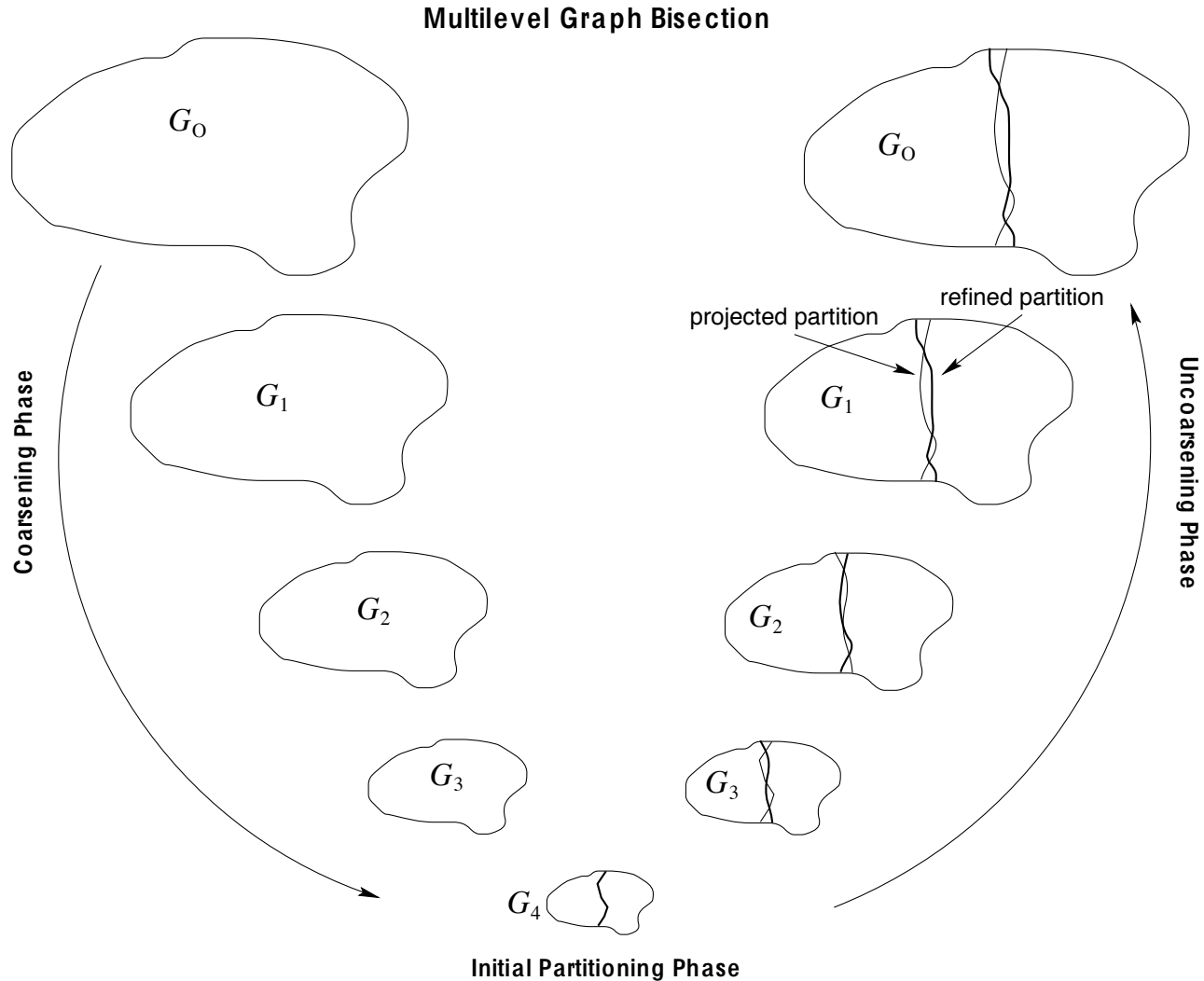
Space-filling curves: Z-Order Curves



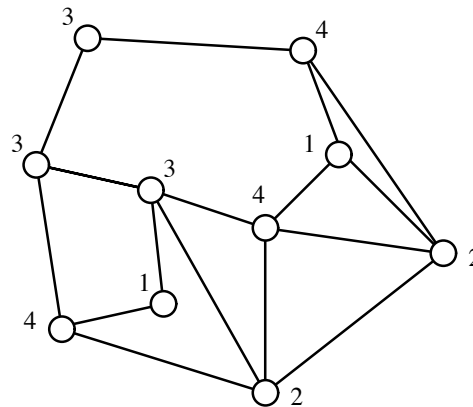
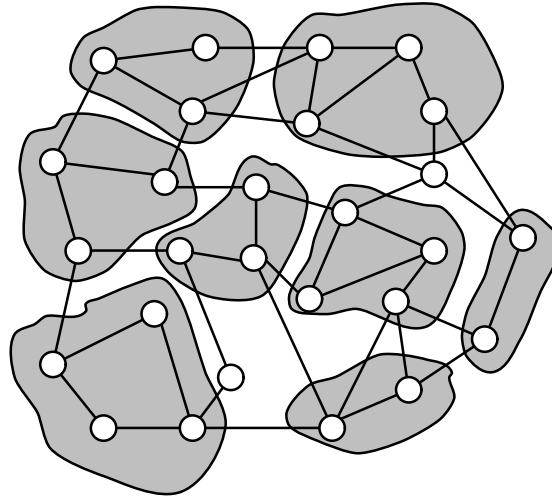
Space-filling curves: Hilbert Curves



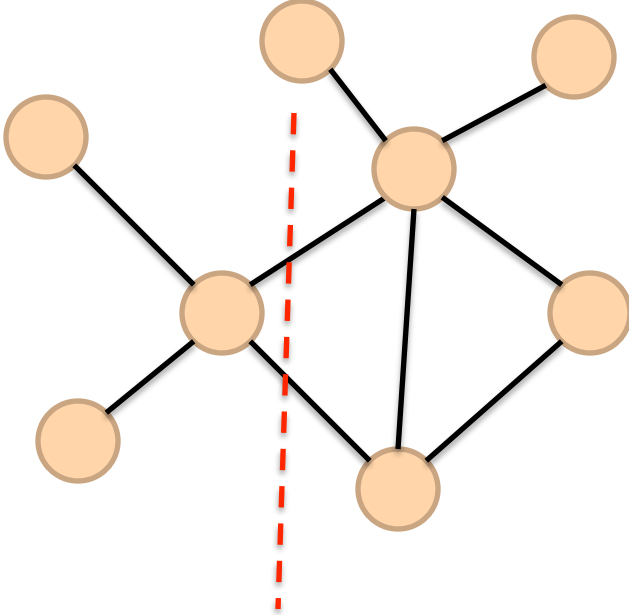
General-Purpose Graph Partitioning



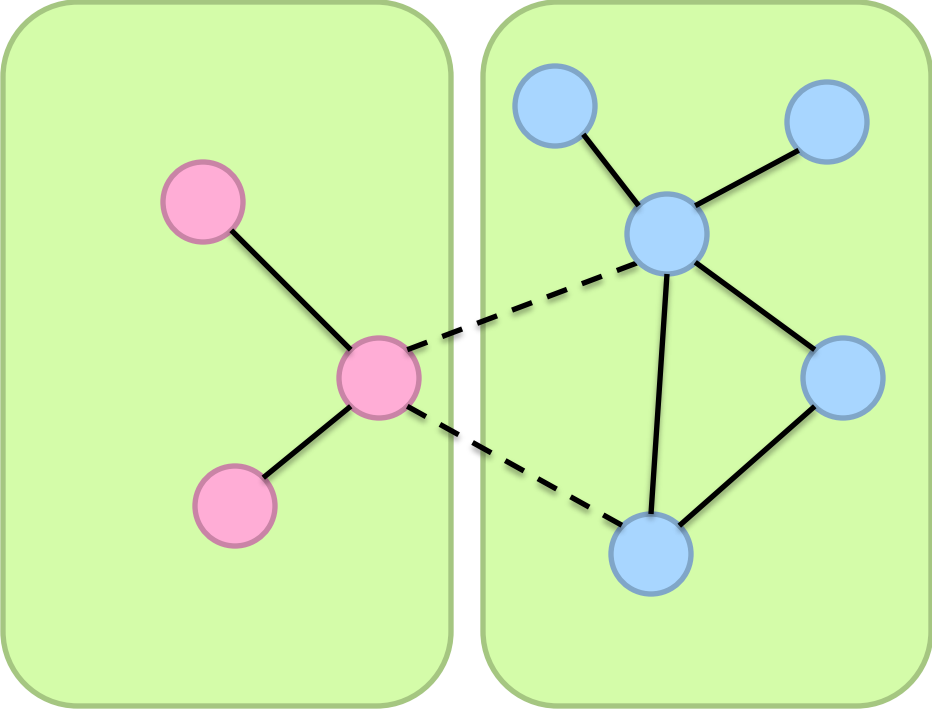
Graph Coarsening



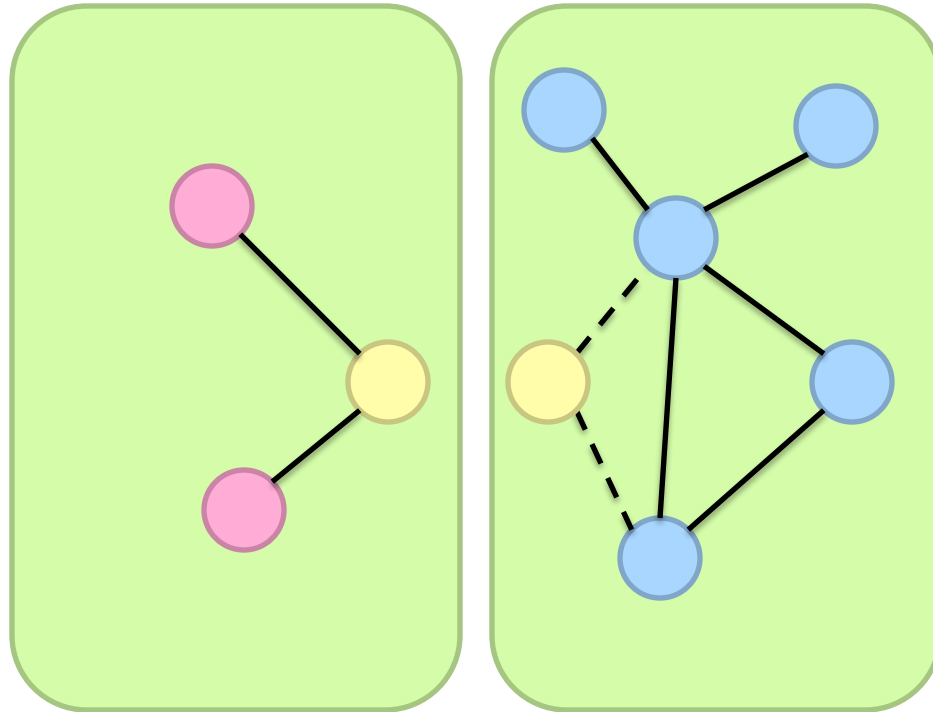
Partition



Partition



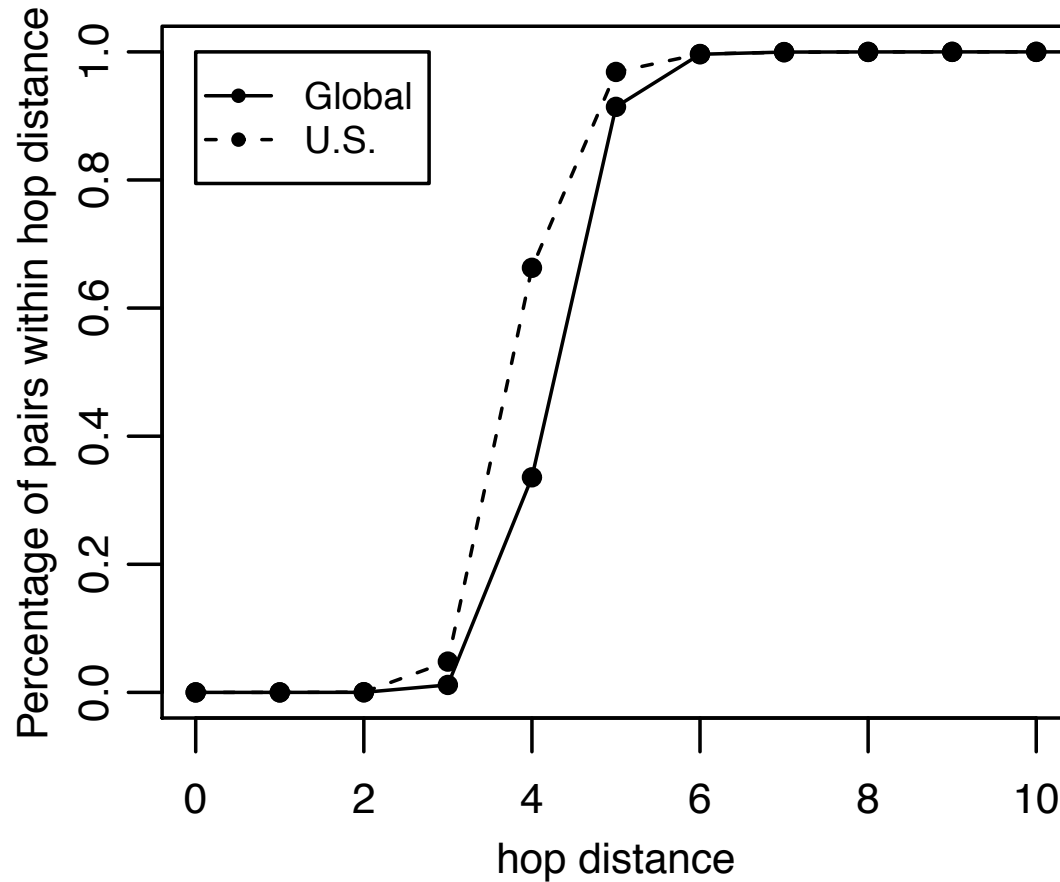
Partition + Replicate



What's the issue?
Solutions?

Neighborhood Replication

What's the cost of replicating n -hop neighborhoods?



What's the more general challenge?

What makes graph processing hard?

- It's tough to apply our “usual tricks” :
 - Partition
 - Replicate
 - Reduce cross-partition communication

Graph Processing Frameworks

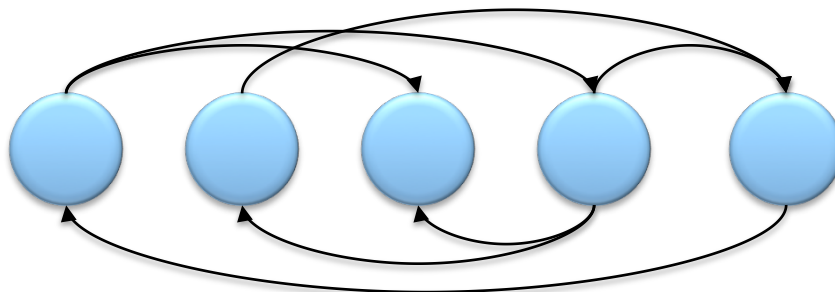


Pregel: Computational Model

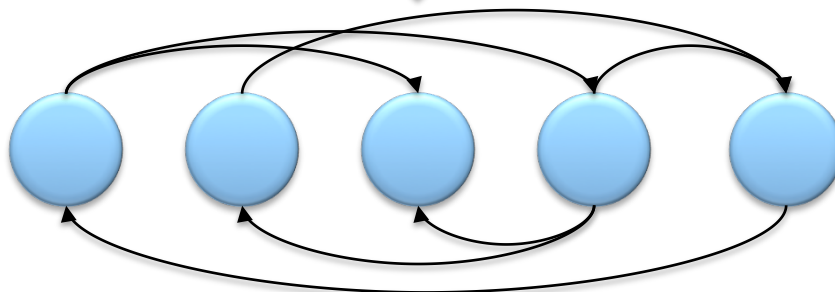
- Based on Bulk Synchronous Parallel (BSP)
 - Computational units encoded in a directed graph
 - Computation proceeds in a series of supersteps
 - Message passing architecture
- Each vertex, at each superstep:
 - Receives messages directed at it from previous superstep
 - Executes a user-defined function (modifying state)
 - Emits messages to other vertices (for the next superstep)
- Termination:
 - A vertex can choose to deactivate itself
 - Is “woken up” if new messages received
 - Computation halts when all vertices are inactive

Pregel

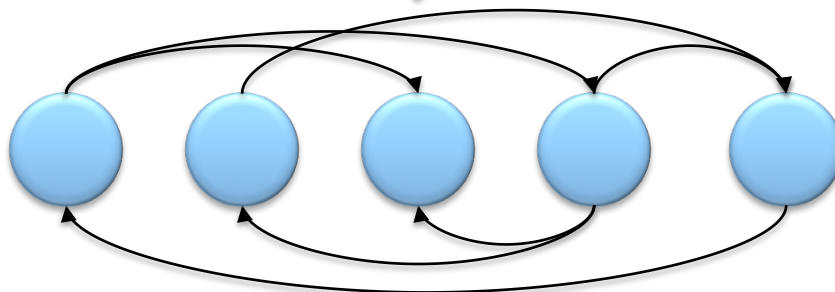
superstep t



superstep $t+1$



superstep $t+2$



Pregel: Implementation

- Master-Slave architecture
 - Vertices are hash partitioned (by default) and assigned to workers
 - Everything happens in memory
- Processing cycle:
 - Master tells all workers to advance a single superstep
 - Worker delivers messages from previous superstep, executing vertex computation
 - Messages sent asynchronously (in batches)
 - Worker notifies master of number of active vertices
- Fault tolerance
 - Checkpointing
 - Heartbeat/revert

Pregel: PageRank

```
class PageRankVertex : public Vertex<double, void, double> {
public:
    virtual void Compute(MessageIterator* msgs) {
        if (superstep() >= 1) {
            double sum = 0;
            for (; !msgs->Done(); msgs->Next())
                sum += msgs->Value();
            *MutableValue() = 0.15 / NumVertices() + 0.85 * sum;
        }

        if (superstep() < 30) {
            const int64 n = GetOutEdgeIterator().size();
            SendMessageToAllNeighbors(GetValue() / n);
        } else {
            VoteToHalt();
        }
    }
};
```

Pregel: SSSP

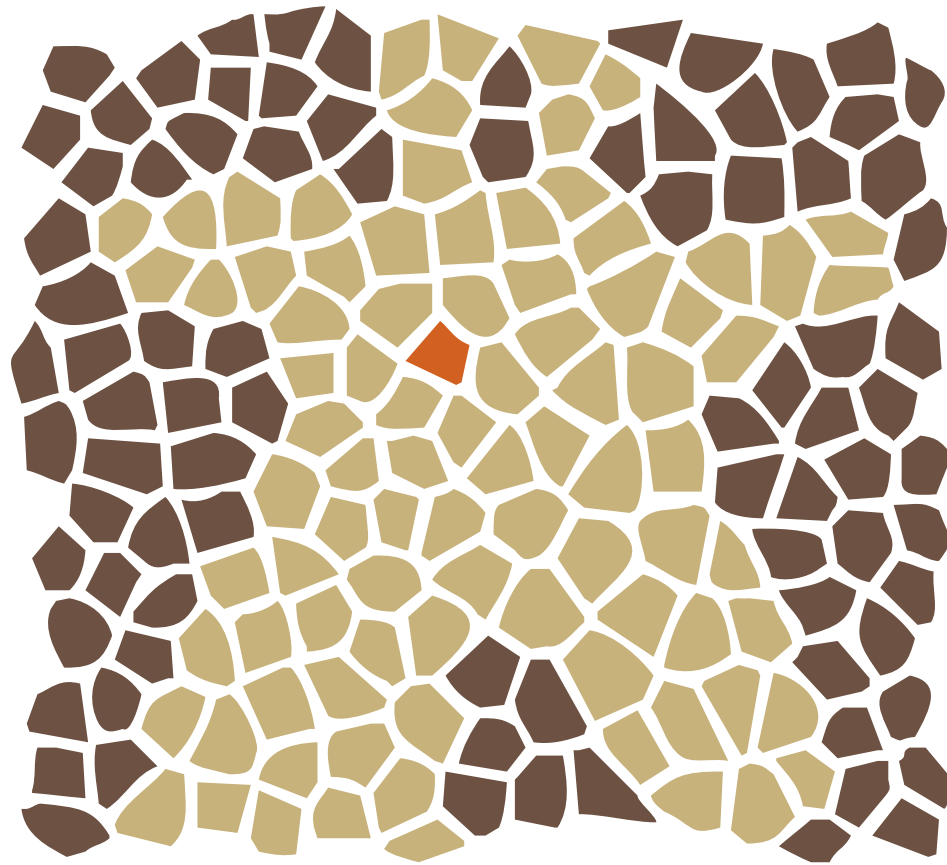
```
class ShortestPathVertex : public Vertex<int, int, int> {
    void Compute(MessageIterator* msgs) {
        int mindist = IsSource(vertex_id()) ? 0 : INF;
        for (; !msgs->Done(); msgs->Next())
            mindist = min(mindist, msgs->Value());
        if (mindist < GetValue()) {
            *MutableValue() = mindist;
            OutEdgeIterator iter = GetOutEdgeIterator();
            for (; !iter.Done(); iter.Next())
                SendMessageTo(iter.Target(),
                    mindist + iter.GetValue());
        }
        VoteToHalt();
    }
};
```


Pregel: Combiners

```
class MinIntCombiner : public Combiner<int> {
    virtual void Combine(MessageIterator* msgs) {

        int mindist = INF;
        for (; !msgs->Done(); msgs->Next())
            mindist = min(mindist, msgs->Value());
        Output("combined_source", mindist);
    }

};
```



A P A C H E
G I R A P H

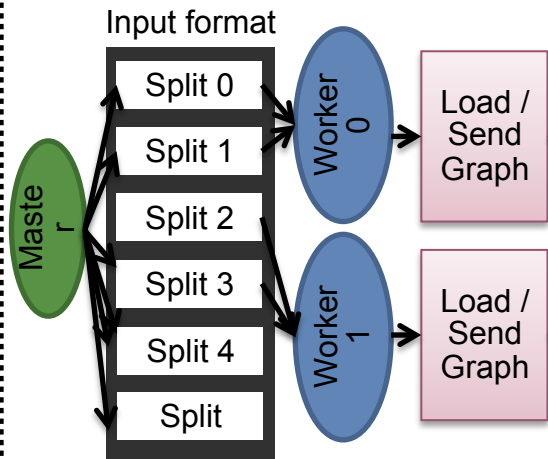
Giraph Architecture

- Master – Application coordinator
 - Synchronizes supersteps
 - Assigns partitions to workers before superstep begins
- Workers – Computation & messaging
 - Handle I/O – reading and writing the graph
 - Computation/messaging of assigned partitions
- ZooKeeper
 - Maintains global application state

Giraph Dataflow

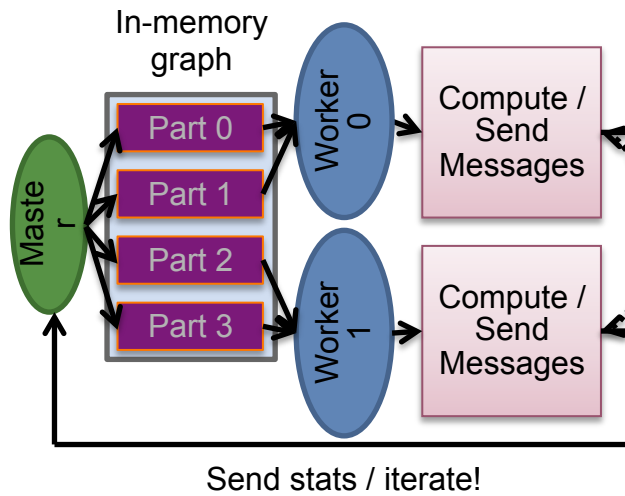
1

Loading the graph



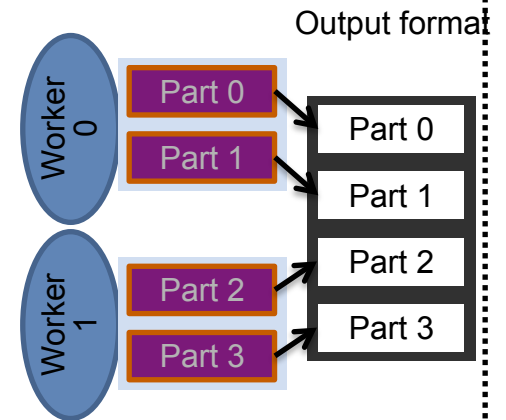
2

Compute/Iterate

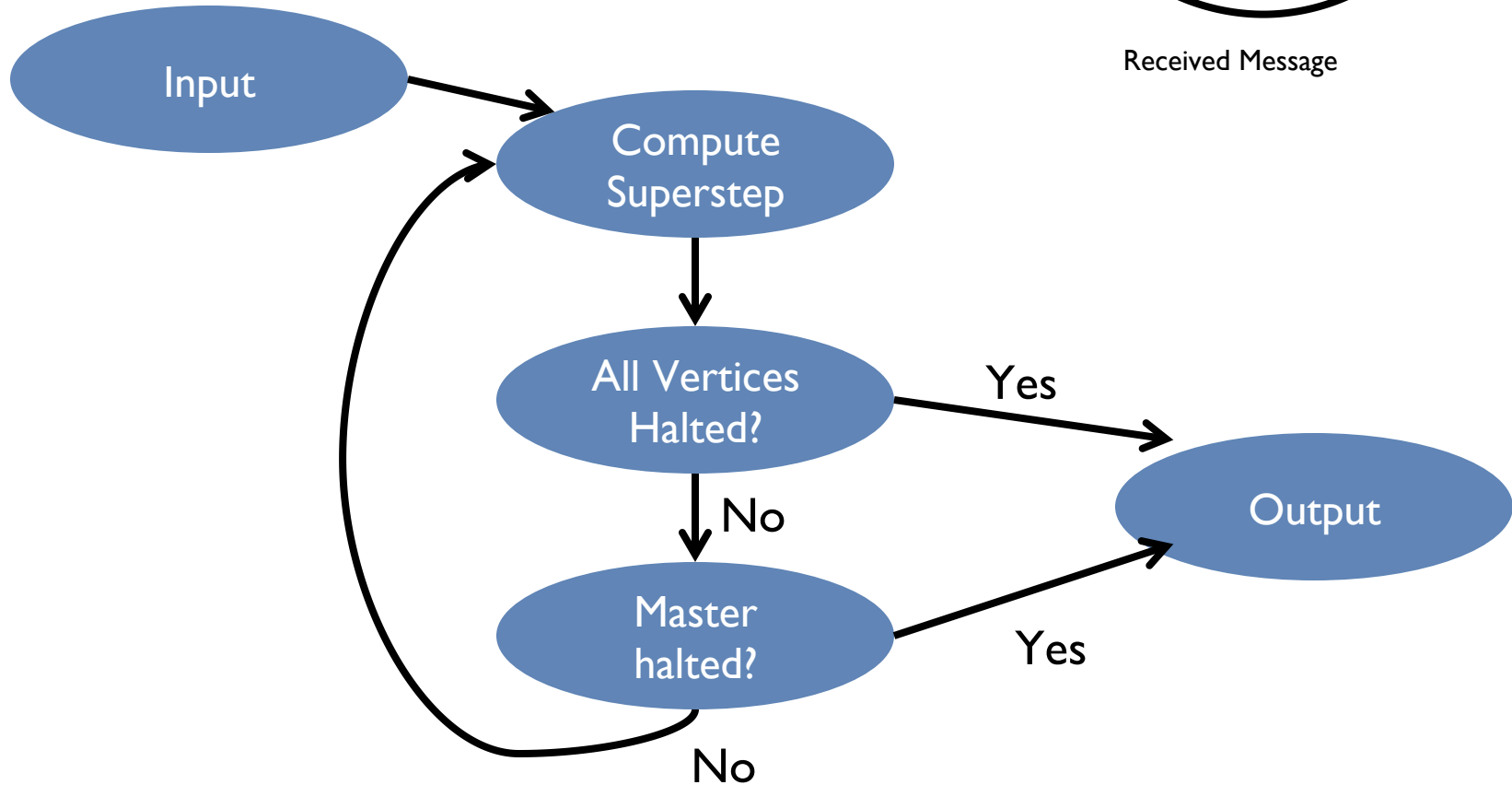


3

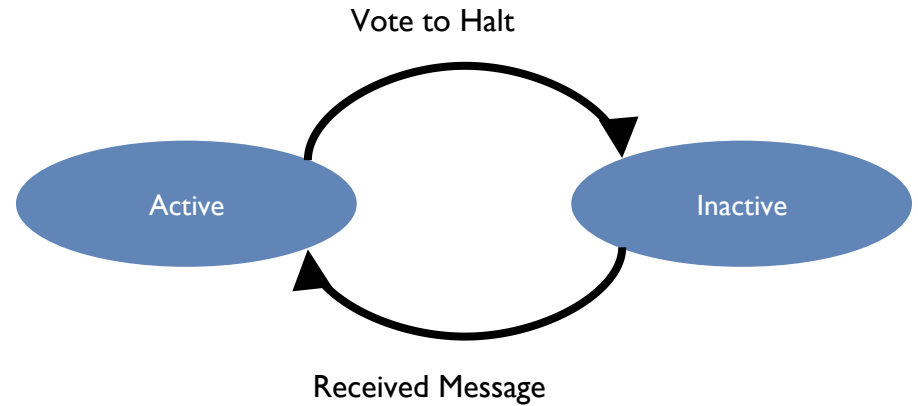
Storing the graph



Giraph Lifecycle



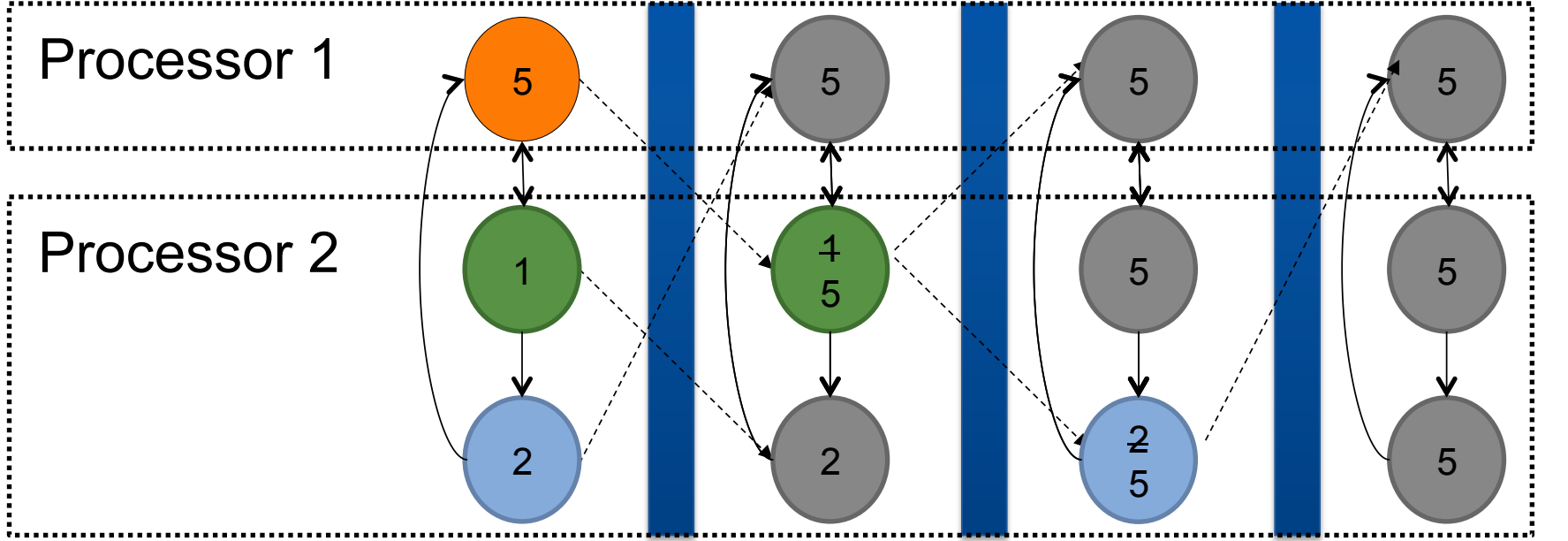
Vertex Lifecycle



Giraph Example

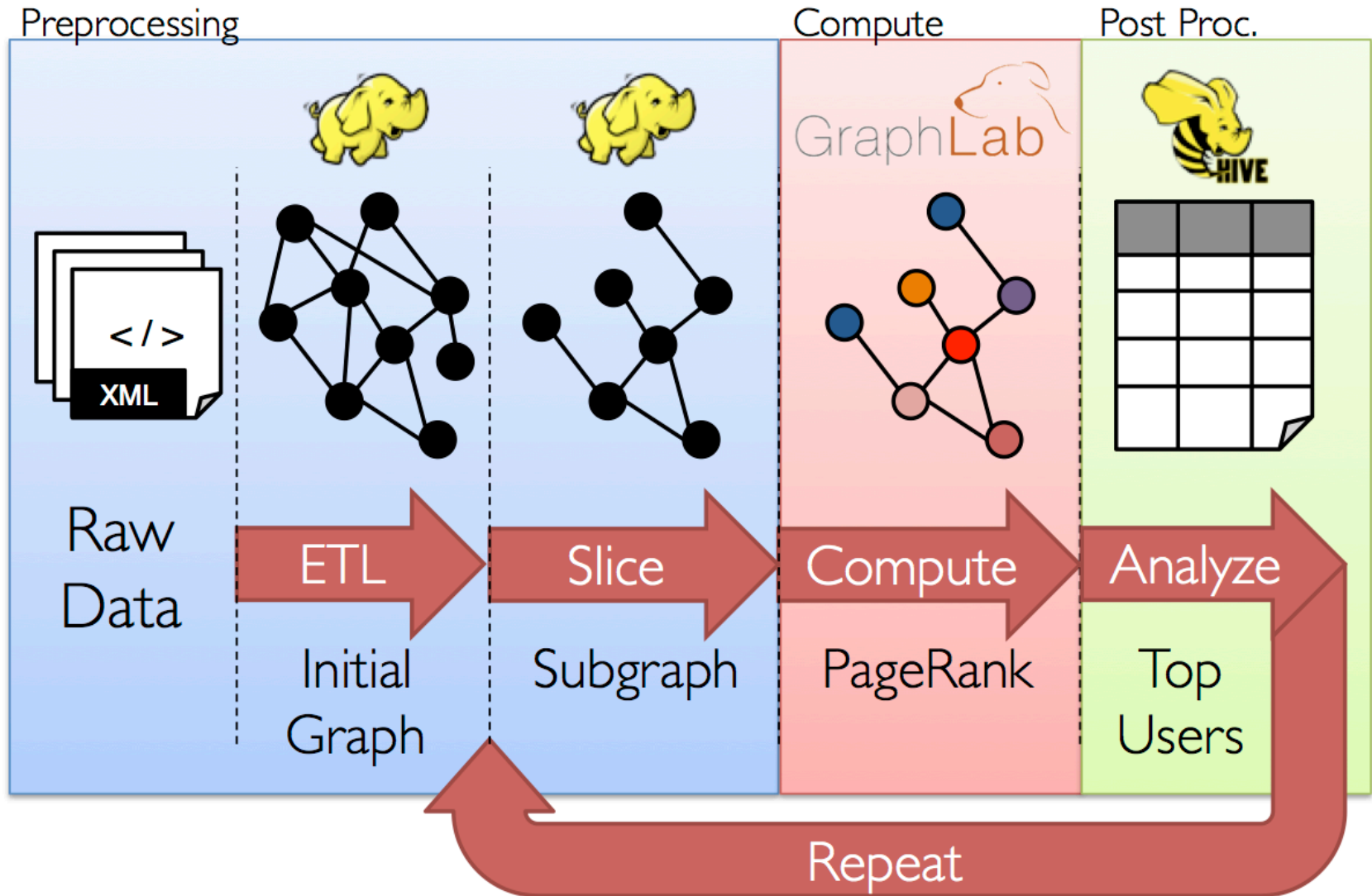
```
public class MaxComputation extends BasicComputation<IntWritable, IntWritable,
    NullWritable, IntWritable> {
    @Override
    public void compute(Vertex<IntWritable, IntWritable, NullWritable> vertex,
        Iterable<IntWritable> messages) throws IOException
    {
        boolean changed = false;
        for (IntWritable message : messages) {
            if (vertex.getValue().get() < message.get()) {
                vertex.setValue(message);
                changed = true;
            }
        }
        if (getSuperstep() == 0 || changed) {
            sendMessageToAllEdges(vertex, vertex.getValue());
        }
        vertex.voteToHalt();
    }
}
```

Execution Trace



Time

GraphX: Motivation

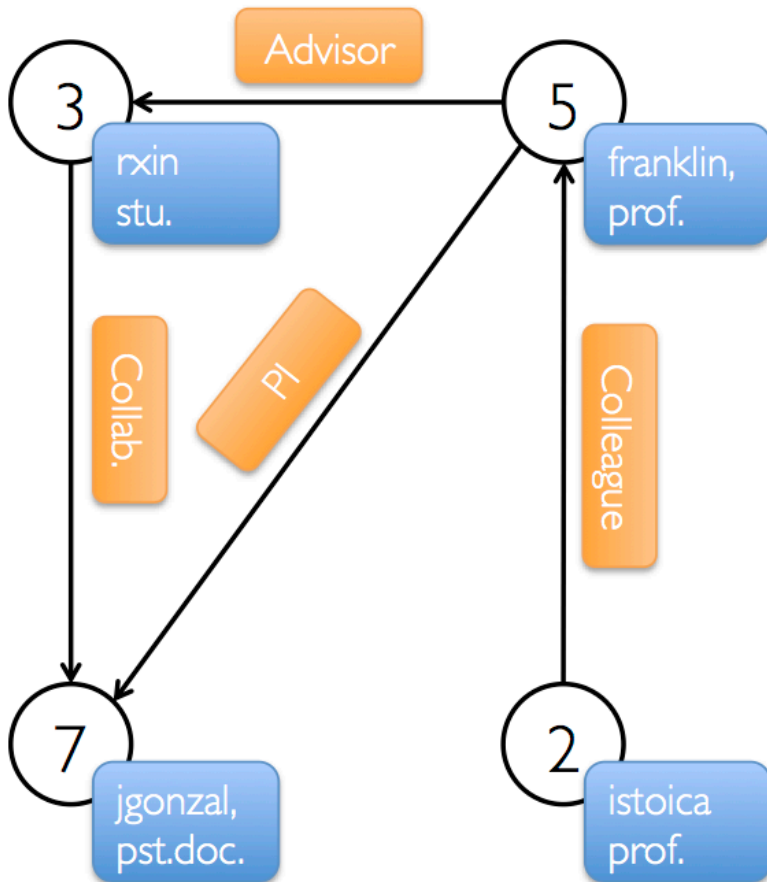


GraphX = Spark for Graphs

- Integration of record-oriented and graph-oriented processing
- Extends RDDs to Resilient Distributed Property Graphs
- Property graphs:
 - Present different views of the graph (vertices, edges, triplets)
 - Support map-like operations
 - Support distributed Pregel-like aggregations

Property Graph: Example

Property Graph



Vertex Table

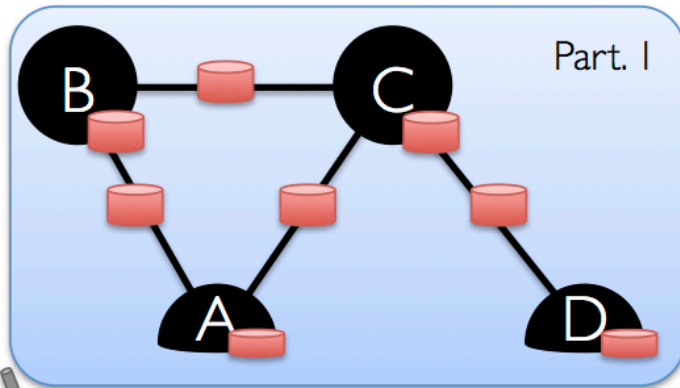
Id	Property (V)
3	(rxin, student)
7	(jgonzal, postdoc)
5	(franklin, professor)
2	(istoica, professor)

Edge Table

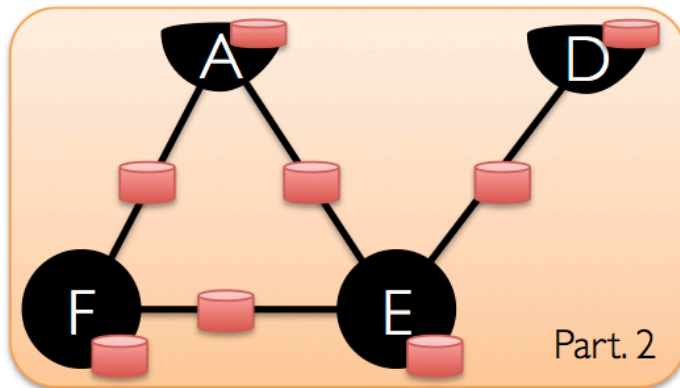
Srclid	Dstld	Property (E)
3	7	Collaborator
5	3	Advisor
2	5	Colleague
5	7	PI

Underneath the Covers

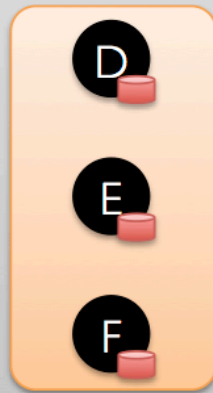
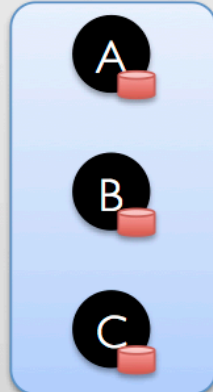
Property Graph



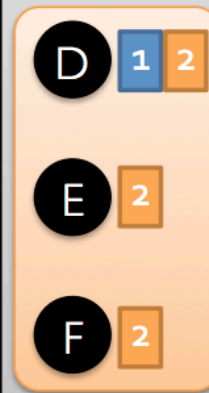
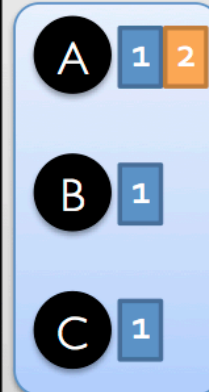
2D Vertex Cut Heuristic



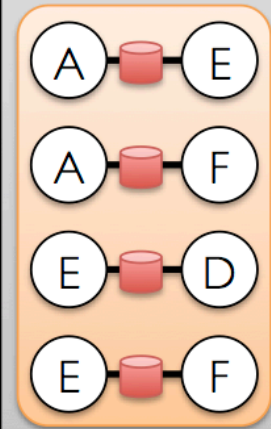
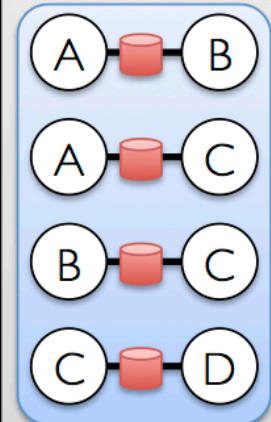
Vertex Table
(RDD)



Routing Table
(RDD)



Edge Table
(RDD)



Today's Agenda

- What makes graph processing hard?
- Graph processing frameworks
- Twitter case study



Questions?