INFM 603: Information Technology and Organizational Context

# **Session 6: Relational Databases**



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# Databases Today...





#### What's structured information? It's what you put in a database

#### What's a database?

It's what you store structured information in

#### So what's a database?

An integrated collection of data organized according to some model...

#### So what's a relational database?

An integrated collection of data organized according to a relational model

#### Database Management System (DBMS)

Software system designed to store, manage, and facilitate access to databases

# Databases (try to) model reality...

- Entities: things in the world
  - Example: airlines, tickets, passengers
- Relationships: how different things are related
  - Example: the tickets each passenger bought
- "Business Logic": rules about the world
  - Example: fare rules



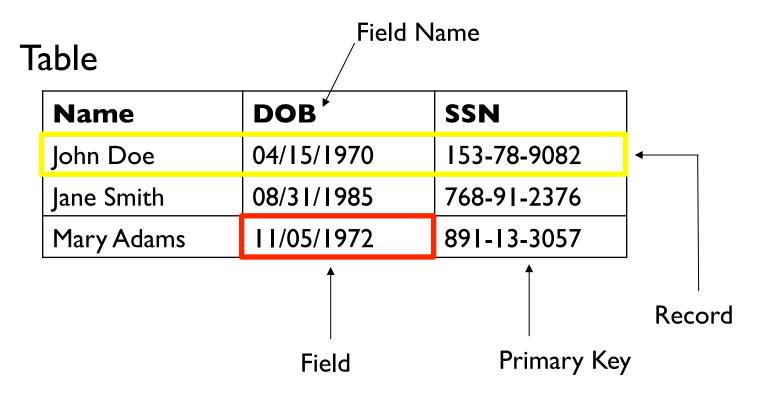
# **Relational Databases**

Source: Microsoft Office Clip Art

### **Components of a Relational Database**

- Field: an "atomic" unit of data
- Record: a collection of related fields
  - Sometimes called a "tuple"
- Table: a collection of related records
  - Each record is a row in the table
  - Each field is a column in the table
- Database: a collection of tables

# A Simple Example



# Why "Relational"?

• View of the world in terms of entities and relations:

- Tables represent "relations"
- Each row (record, tuple) is "about" an entity
- Fields can be interpreted as "attributes" or "properties" of the entity
- Data is manipulated by "relational algebra":
  - Defines things you can do with tuples
  - Expressed in SQL

### The Registrar Example

- What do we need to know?
  - Something about the students (e.g., first name, last name, email, department)
  - Something about the courses (e.g., course ID, description, enrolled students, grades)
  - Which students are in which courses
- How do we capture these things?

## A First Try

Student ID	Last Name	First Name	Dept ID	Dept	Course ID	Course name	Grade	email
1	Arrows	John	EE	EE	lbsc690	Information Technology	90	jarrows@wam
1	Arrows	John	EE	Elec Engin	ee750	Communication	95	ja 2002@yahoo
2	Peters	Kathy	HIST	HIST	lbsc690	Informatino Technology	95	kpeters2@wam
2	Peters	Kathy	HIST	history	hist405	American History	80	kpeters2@wma
3	Smith	Chris	HIST	history	hist405	American History	90	smith2002@glue
4	Smith	John	CLIS	Info Sci	lbsc690	Information Technology	98	js03@wam

#### Why is this a bad idea?

# **Goals of "Normalization"**

- Save space
  - Save each fact only once
- More rapid updates
  - Every fact only needs to be updated once
- More rapid search
  - Finding something once is good enough
- Avoid inconsistency
  - Changing data once changes it everywhere

### Another Try...

#### Student Table

Student ID	Last Name	First Name	Department ID	email
1	Arrows	John	EE	jarrows@wam
2	Peters	Kathy	HIST	kpeters2@wam
3	Smith	Chris	HIST	smith2002@glue
4	Smith	John	CLIS	<u>js03@wam</u>

#### Department Table

Department ID	Department
EE	Electrical Engineering
HIST	History
CLIS	Information Studies

#### Course Table

Course ID	Course Name
lbsc690	Information Technology
ee750	Communication
hist405	American History

#### **Enrollment Table**

Student ID	Course ID	Grade
1	lbsc690	90
1	ee750	95
2	lbsc690	95
2	hist405	80
3	hist405	90
4	lbsc690	98

# Keys

- "Primary Key" uniquely identifies a record
  - e.g., student ID in the student table
- "Foreign Key" is primary key in the other table
  - It need not be unique in this table



# **Approaches to Normalization**

- For simple problems:
  - Start with the entities you're trying to model
  - Group together fields that "belong together"
  - Add keys where necessary to connect entities in different tables
- For more complicated problems:
  - Entity-relationship modeling

### The Data Model

#### Student Table

Student ID	Last Name	First Name	Department ID	email
1	Arrows	John	EE	jarrows@wam
2	Peters	Kathy	HIST	kpeters2@wam
3	Smith	Chris	HIST	smith2002@glue
4	Smith	John	CLIS	<u>js03@wam</u>

#### Department Table

Department ID	Department
EE	Electrical Engineering
HIST	History
CLIS	Information Studies

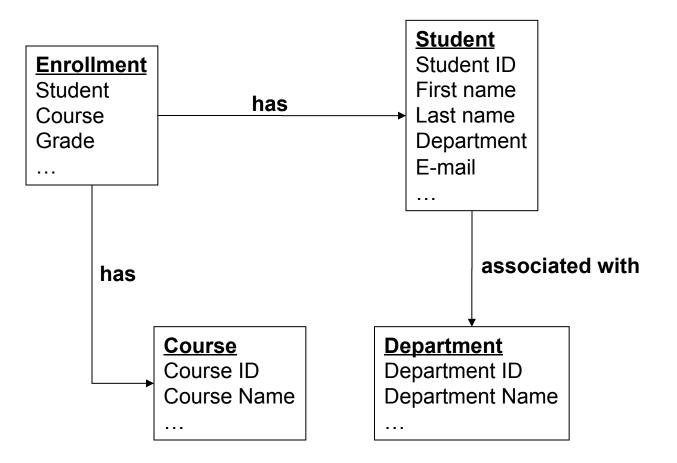
#### Course Table

Course ID	Course Name
lbsc690	Information Technology
ee750	Communication
hist405	American History

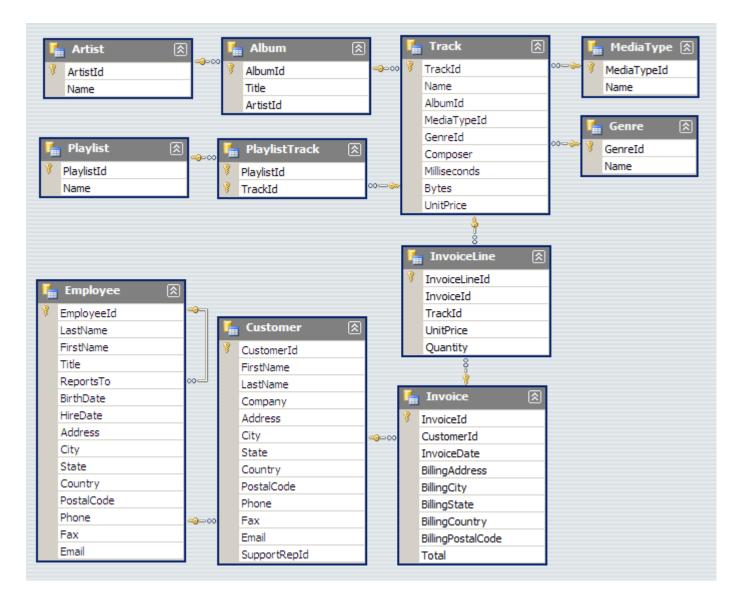
#### **Enrollment Table**

Student ID	Course ID	Grade
1	lbsc690	90
1	ee750	95
2	lbsc690	95
2	hist405	80
3	hist405	90
4	lbsc690	98

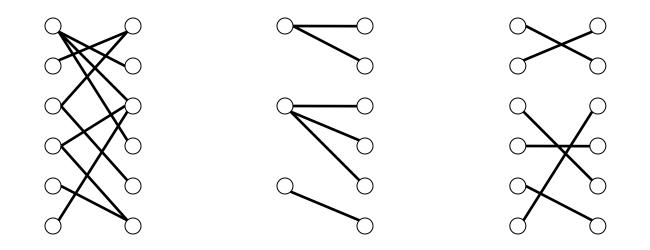
### **Registrar ER Diagram**



# **A Real Example**



#### **Types of Relationships**



Many-to-Many

One-to-Many

One-to-One

# **Database Integrity**

- Registrar database must be internally consistent
  - All enrolled students must have an entry in the student table
  - All courses must have a name
  - ...
- What happens:
  - When a student withdraws from the university?
  - When a course is taken off the books?

# **Integrity Constraints**

• Conditions that must be true of the database at any time

- Specified when the database is designed
- Checked when the database is modified
- RDBMS ensures that integrity constraints are always kept
  - So that database contents remain faithful to the real world
  - Helps avoid data entry errors
- Where do integrity constraints come from?



#### Select

Student ID	Last Name	First Name	Dept ID	Department	email
1	Arrows	John	EE	Electrical Engineering	jarrows@wam
2	Peters	Kathy	HIST	History	kpeters2@wam
3	Smith	Chris	HIST	History	smith2002@glue
4	Smith	John	CLIS	Information Stuides	<u>js03@wam</u>

#### select Student ID, Department

Student ID	Department
1	Electrical Engineering
2	History
3	History
4	Information Stuides

#### Where

Student ID	Last Name	First Name	Dept ID	Department	email
1	Arrows	John	EE	Electrical Engineering	jarrows@wam
2	Peters	Kathy	HIST	History	kpeters2@wam
3	Smith	Chris	HIST	History	smith2002@glue
4	Smith	John	CLIS	Information Stuides	<u>js03@wam</u>

#### where Department ID = "HIST"

Student ID	Last Name	First Name	Department ID	Department	email
2	Peters	Kathy	HIST	History	kpeters2@wam
3	Smith	Chris	HIST	History	smith2002@glue

# Simple SQL Statements

- Choosing columns: SELECT
  - Based on their labels (field names)
  - \* is a shorthand for saying "all fields"
- Choosing rows: WHERE
  - Based on their contents
     department ID = "HIST"
- These can be specified together

# Simple SQL Template

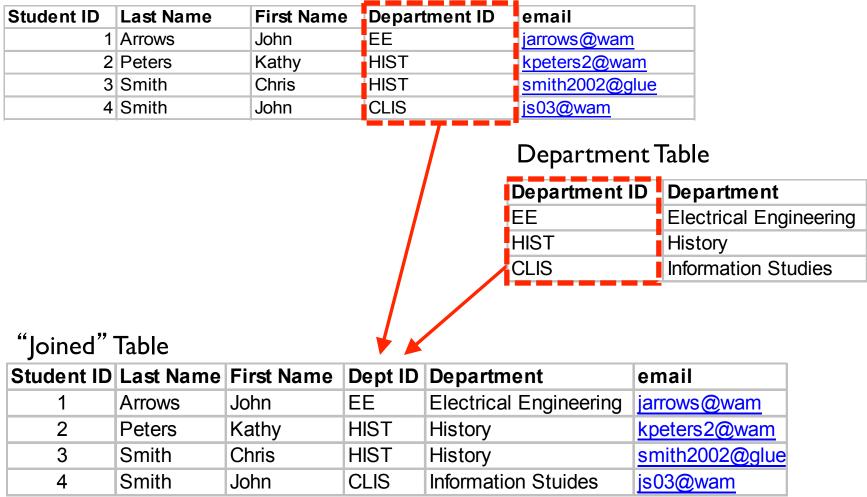
select [columns in the table]
from [table name]
where [selection criteria]

# **SQL** Tips and Tricks

- Referring to fields (in SELECT statements)
  - Use TableName.FieldName
  - Can drop TableName if FieldName is unambiguous
- Selection criteria
  - Use = instead of ==
- Note, different dialects of SQL!



#### Student Table



# **SQL** Template for Joins

select [columns in the table]
from [table name]
join [another tablename] on [join criterion]
join [another tablename] on [join criterion]
...

where [selection criteria]

Join criterion: usually, based on primary/foreign key relationships e.g., Table I.PrimaryKey = Table2.ForeignKey

# Aggregations

• SQL aggregation functions

- Examples: count, min, max, sum, avg
- Use in select statements

```
select count(*)...
select min(price)...
select sum(length)...
```

- Tip: when trying to write SQL query with aggregation, do it first without
- Group by [field]
  - Often used in conjunction with aggregation
  - Conceptually, breaks table apart based on the [field]

### How do you want your results served?

- Order by [field name]
  - Does exactly what you think it does!
  - Either "asc" or "desc"
- Limit *n* 
  - Returns only *n* records
  - Useful to retrieving the top *n* or bottom *n*

#### So how's a database more than a spreadsheet?

# Database in the "Real World"

• Typical database applications:

- Banking (e.g., saving/checking accounts)
- Trading (e.g., stocks)
- Traveling (e.g., airline reservations)
- Social media (e.g., Facebook)
- ...

#### • Characteristics:

- Lots of data
- Lots of concurrent operations
- Must be fast
- "Mission critical" (well... sometimes)

# **Operational Requirements**

- Must hold a lot of data
- Must be reliable
- Must be fast
- Must support concurrent operations

#### Must hold a lot of data Solution: Use lots of machines (Each machine holds a small slice)

So which machine has your copy?

#### Must be reliable Solution: Use lots of machines (Store multiple copies)

But which copy is the right one? How do you keep the copies in sync?

#### Must be fast Solution: Use lots of machines (Share the load)

How do you spread the load?

# Must support concurrent operations

#### Solution: this is hard!

(But fortunately doesn't matter for many applications)

### **Database Transactions**

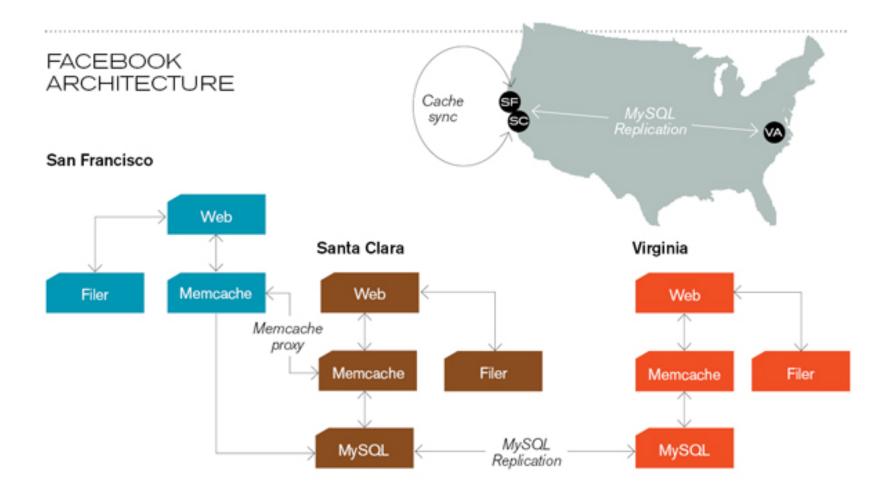
- Transaction = sequence of database actions grouped together
  - e.g., transfer \$500 from checking to savings
- ACID properties:
  - Atomicity: all-or-nothing
  - Consistency: each transaction yield a consistent state
  - Isolation: concurrent transactions must appear to run in isolation
  - Durability: results of transactions must survive even if systems crash

# **Making Transactions**

- Idea: keep a log (history) of all actions carried out while executing transactions
  - Before a change is made to the database, the corresponding log entry is forced to a safe location



- Recovering from a crash:
  - Effects of partially executed transactions are undone
  - Effects of committed transactions are redone
  - Trickier than it sounds!



**Caching servers:** 15 million requests per second, 95% handled by memcache (15 TB of RAM)

**Database layer:** 800 eight-core Linux servers running MySQL (40 TB user data)

